SECISDesign - A Method to Design New and Recombinant Selenoproteins

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Outline

- Biological Introduction
- The Computational Problem
- The Algorithm (SECISDesign)
- Results and Summary

Biological Introduction - 1

Selenoproteins:

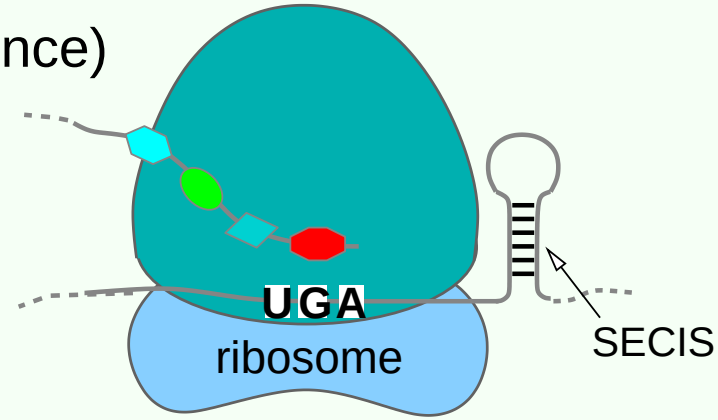
- contain selenium, incorporated as *selenocysteine* at the active site
- greatly enhanced enzymatic activities compared to the cysteine homologues

- important to human health:
  - thyroid hormone metabolism
  - immune function
  - protection against cancer

Biological Introduction - 2

Selenocysteine: (= 21*st* amino acid)

- encoded by the STOP-codon UGA
- inserted, if:
  - UGA is followed by a *SECIS-element*
    (hairpin-like structure + specific sequence)



Biological Introduction - 2

Selenocysteine: (= 21*st* amino acid)

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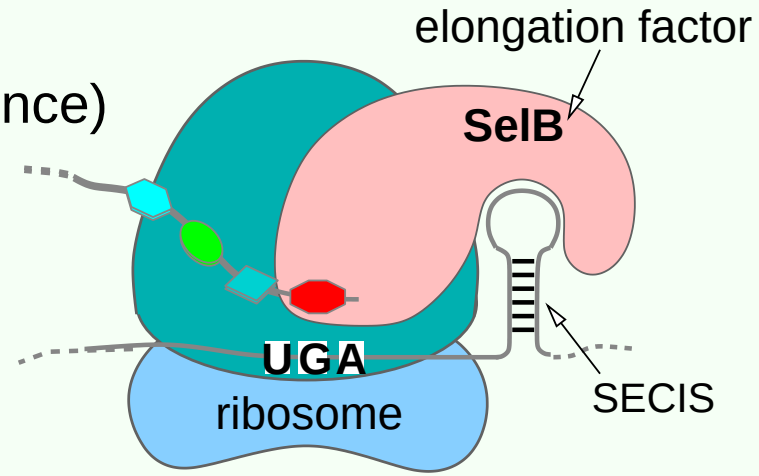
- inserted, if:
  - UGA is followed by a *SECIS-element* (hairpin-like structure + specific sequence)
  - special elongation factor SelB available

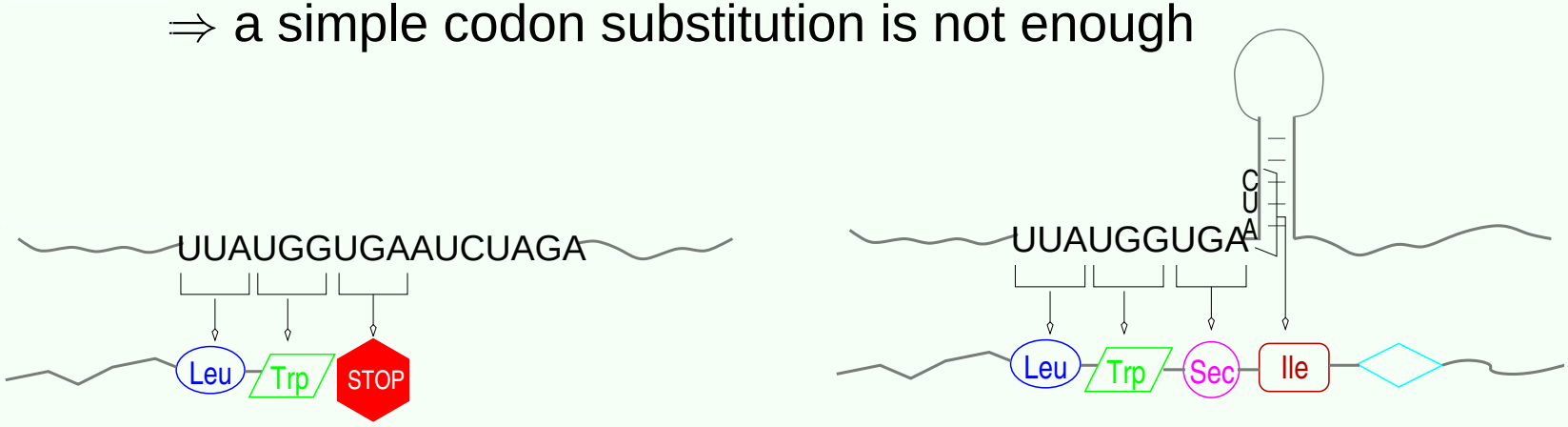


The Biological Problem - 1

Tasks:

1. Replacement of an amino acid (e.g. cysteine) by selenocysteine

$\Rightarrow$ a simple codon substitution is not enough



2. Expression of eukaryotic selenoproteins in *E.coli*

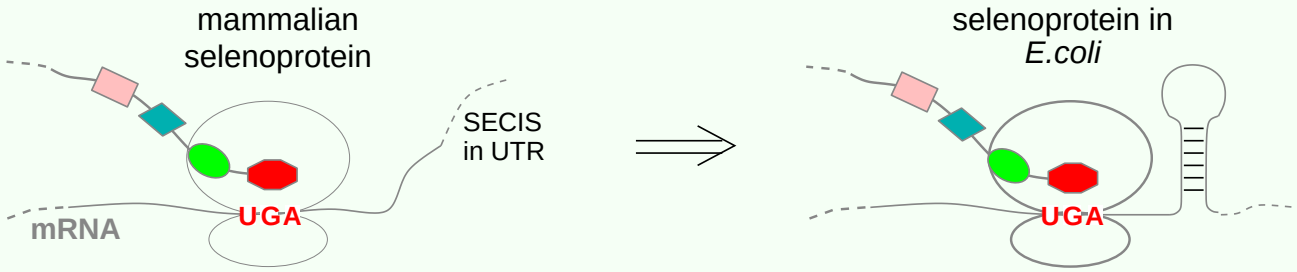
The Biological Problem - 2

Protein expression system: *E.coli*

Eukaryotes: SECIS-element *outside* the coding sequence

E.coli (bacteria): SECIS-element immediately downstream the UGA-Codon

→ located *inside* the coding region of the protein



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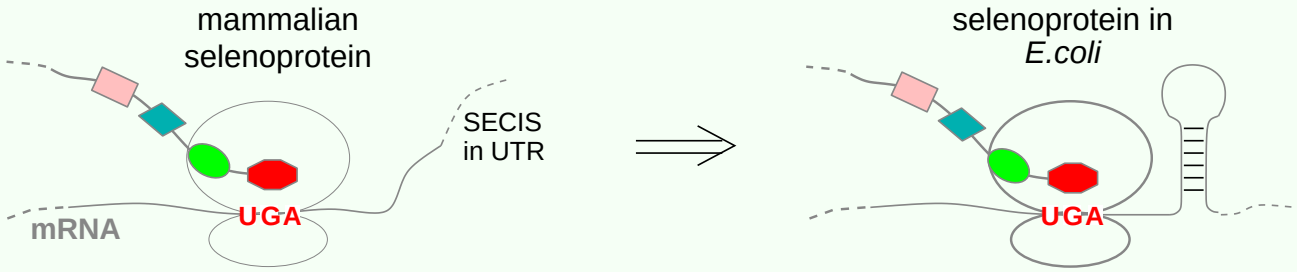
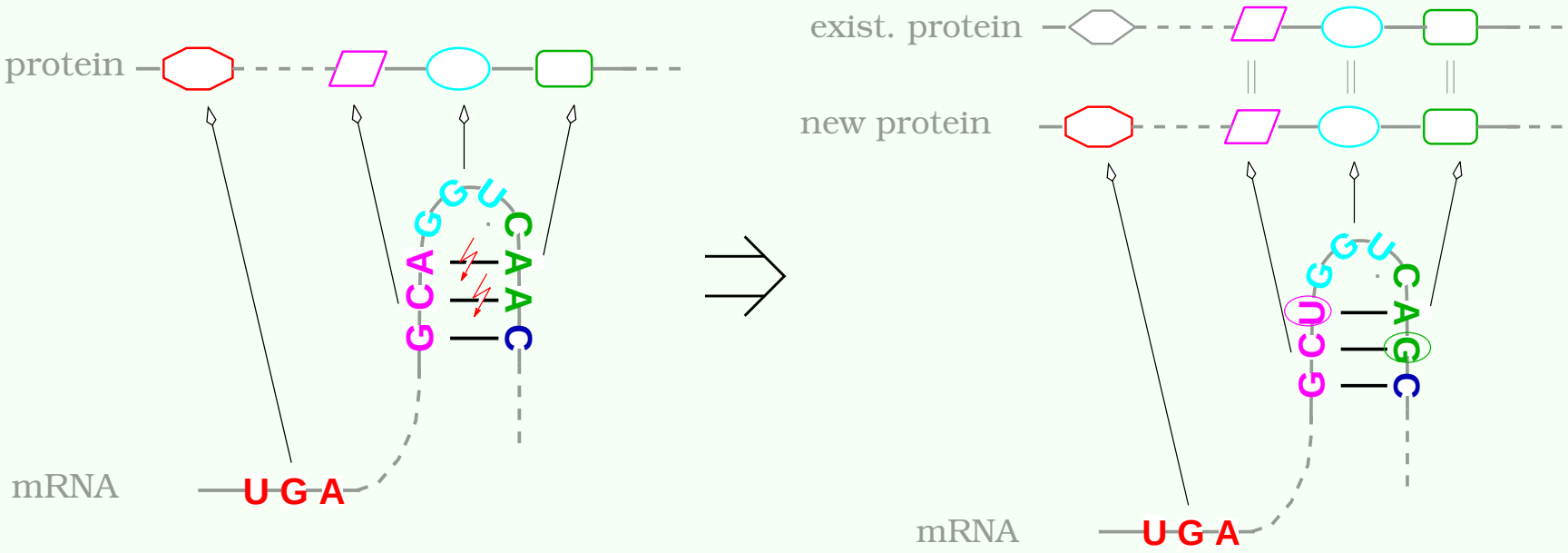


Consequence:

- eukaryotic selenoproteins can't be expressed directly in *E.coli*
- design of a SECIS-element next to the UGA-position
- this design may change the sequence of the protein

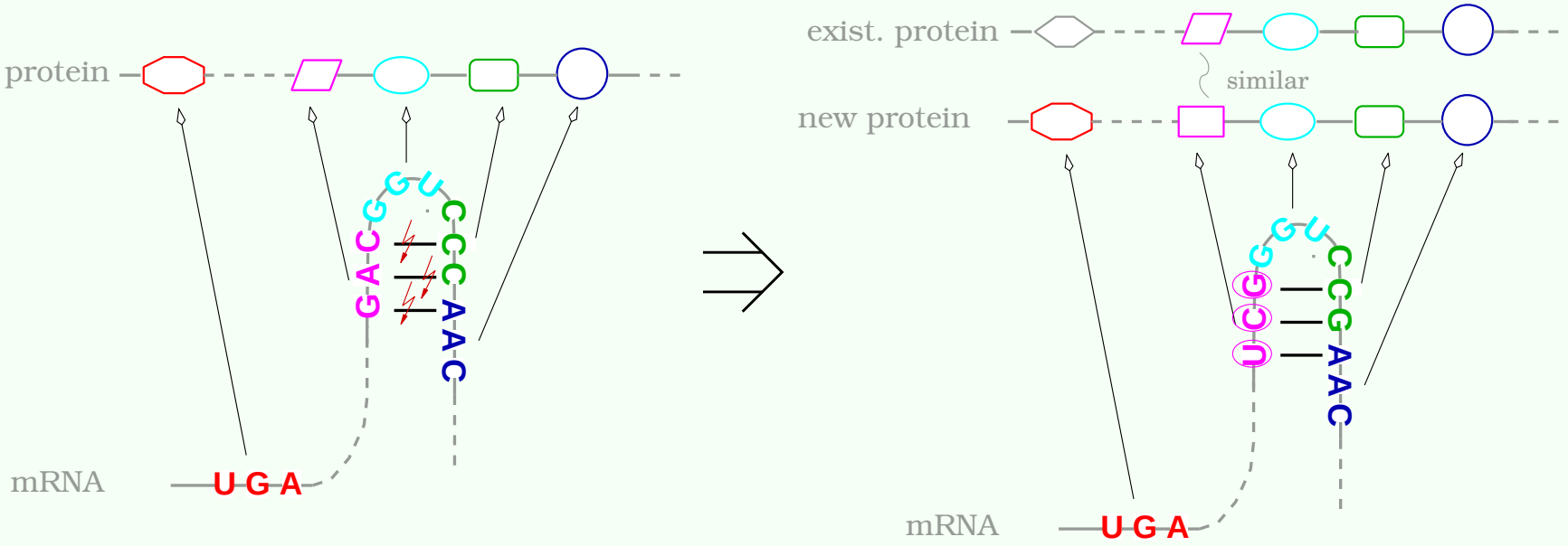
The Biological Problem - 3

SECIS-element-design: (sometimes without a change in the amino acid sequence...)



The Biological Problem - 4

...and sometimes wobble mutations are NOT enough!



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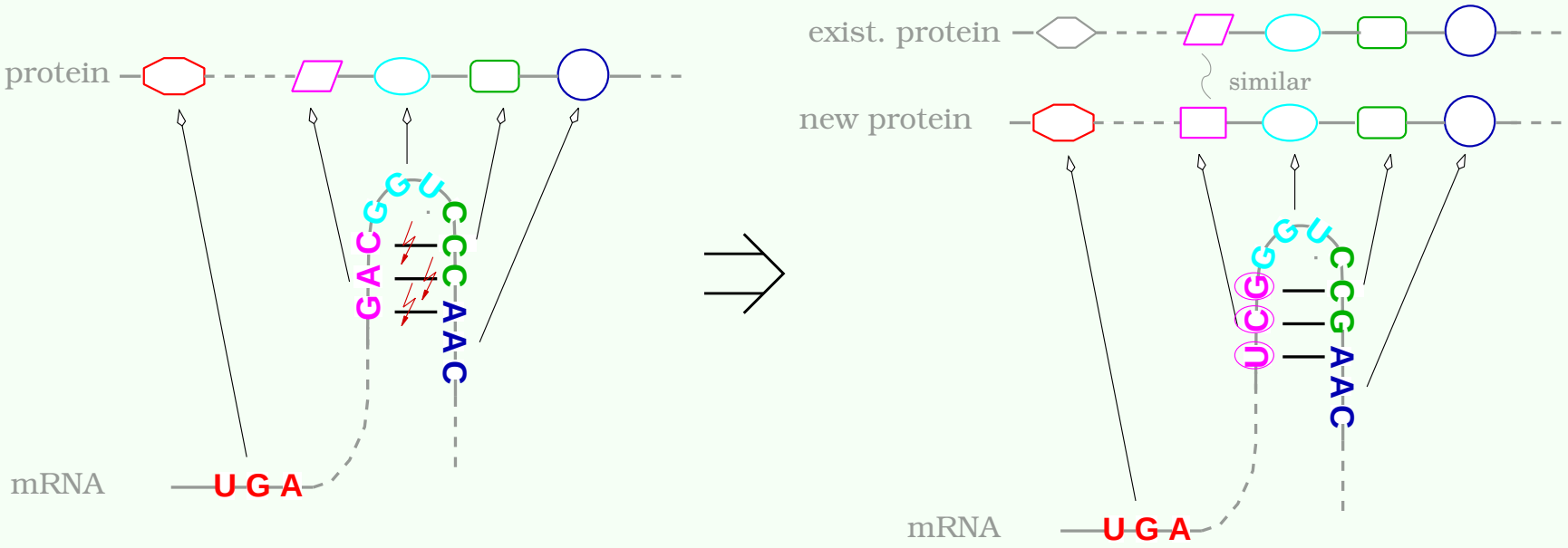
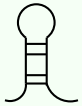


Problem:

compromise between

- quality of the SECIS-element and
- changes in the protein sequence

The Computational Problem

- **Given:** $G =$  ... typical SECIS secondary structure

  $S = S_1 \ldots S_{3n}$ ... nucleotide sequence (SECIS-consensus)

  $A = A_1 \ldots A_n$ ... original amino acid sequence

$$A = A_1 \quad \ldots \quad A_i \quad \ldots \quad A_n$$

$$S = S_1 S_2 S_3 \quad \ldots \quad S_{3i-2} S_{3i-1} S_{3i} \quad \ldots \quad S_{3n-2} S_{3n-1} S_{3n}$$

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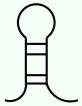
- **Given:** $G =$  ... typical SECIS secondary structure

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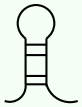
  · can adopt $G =$ 

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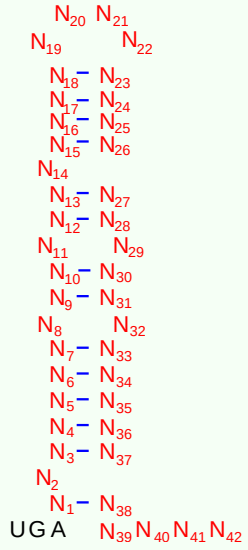

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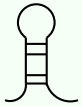
- **Given:** $G =$  ... typical SECIS secondary structure

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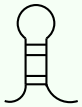
  - can adopt $G =$ 
  - has maximum similarity with S

$$A = A_1 \quad \ldots \quad A_i \quad \ldots \quad A_n$$

$$N = \overbrace{N_1 N_2 N_3} \ \ldots \ \overbrace{N_{3i-2} N_{3i-1} N_{3i}} \ \ldots \ \overbrace{N_{3n-2} N_{3n-1} N_{3n}}$$

$$\sim \sim \sim \qquad \sim \sim \sim \qquad \sim \sim \sim$$

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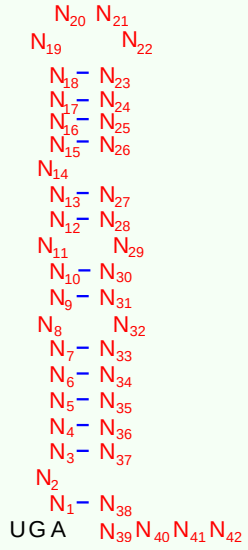

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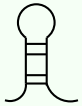
- **Given:** $G =$  ... typical SECIS secondary structure

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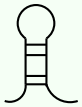
  · can adopt $G =$ 

  · has maximum similarity with S

  · encodes amino acid sequence A' with maximal similarity to A

$$
\begin{array}{lccccccc}
A = & A_1 & \ldots & A_i & \ldots & A_n \\
 & \sim & & \sim & & \sim \\
A' = & A'_1 & \ldots & A'_i & \ldots & A'_n \\
N = & \overbrace{N_1 N_2 N_3} & \ldots & \overbrace{N_{3i-2} N_{3i-1} N_{3i}} & \ldots & \overbrace{N_{3n-2} N_{3n-1} N_{3n}} \\
 & \sim\ \sim\ \sim & & \sim\ \sim\ \sim & & \sim\ \sim\ \sim \\
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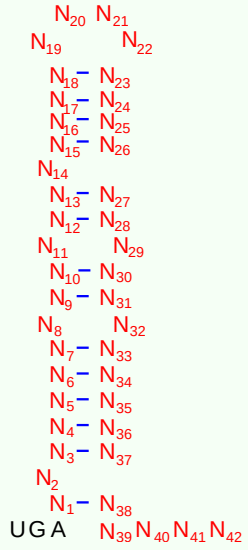
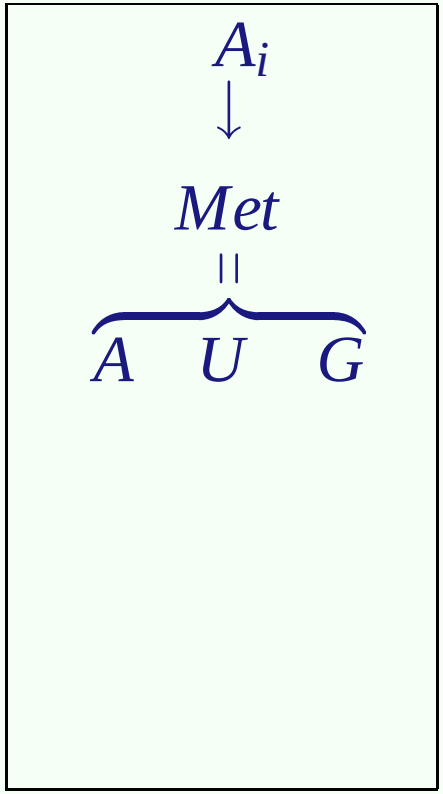
\end{array}
$$



Contrary Conditions

Possible Problem: contrary conditions at nucleotide and amino acid level



→ essential amino acid

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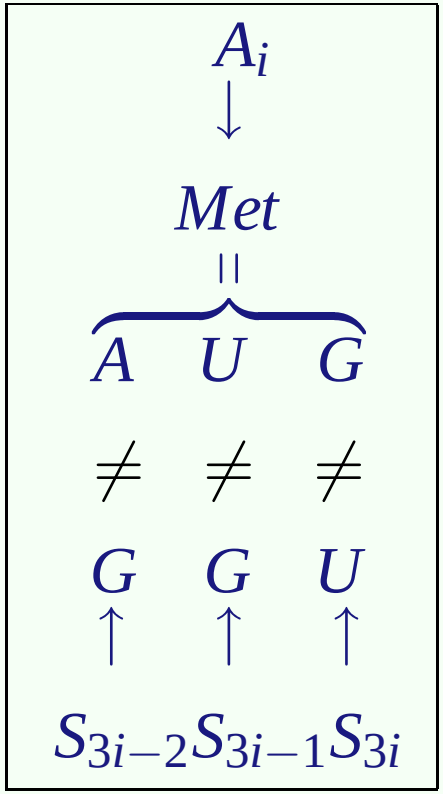

→ essential amino acid

→ essential SECIS-sequence

⇒ insertions/deletions & *optional* bonds

Extentions - 1

(1) Insertions and deletions (amino acids)

- single amino acids are allowed to be deleted or inserted
- mapping between A and A' changes

$A:$	A_1	$-$	A_2
$A':$	A'_1	A'_2	A'_3
$N:$	$N_1N_2N_3$	$N_4N_5N_6$	$N_7N_8N_9$
$S:$	$S_1S_2S_3$	$S_4S_5S_6$	$S_7S_8S_9$
SECpos	1	2	3

- insertion of A'_2
- A'_2 has no counterpart in A
- insertion penalty IP

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$A:$	A_1	A_2	A_3	A_4
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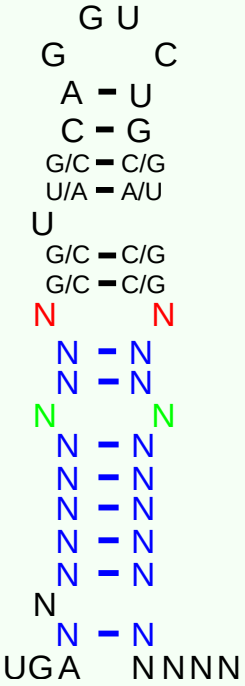
$S:$	$S_1S_2S_3$	$---$	$S_4S_5S_6$	$S_7S_8S_9$
SECpos	1	2		3

- deletion of A_2
- compare $N_4N_5N_6$ with A_3
- deletion penalty DP

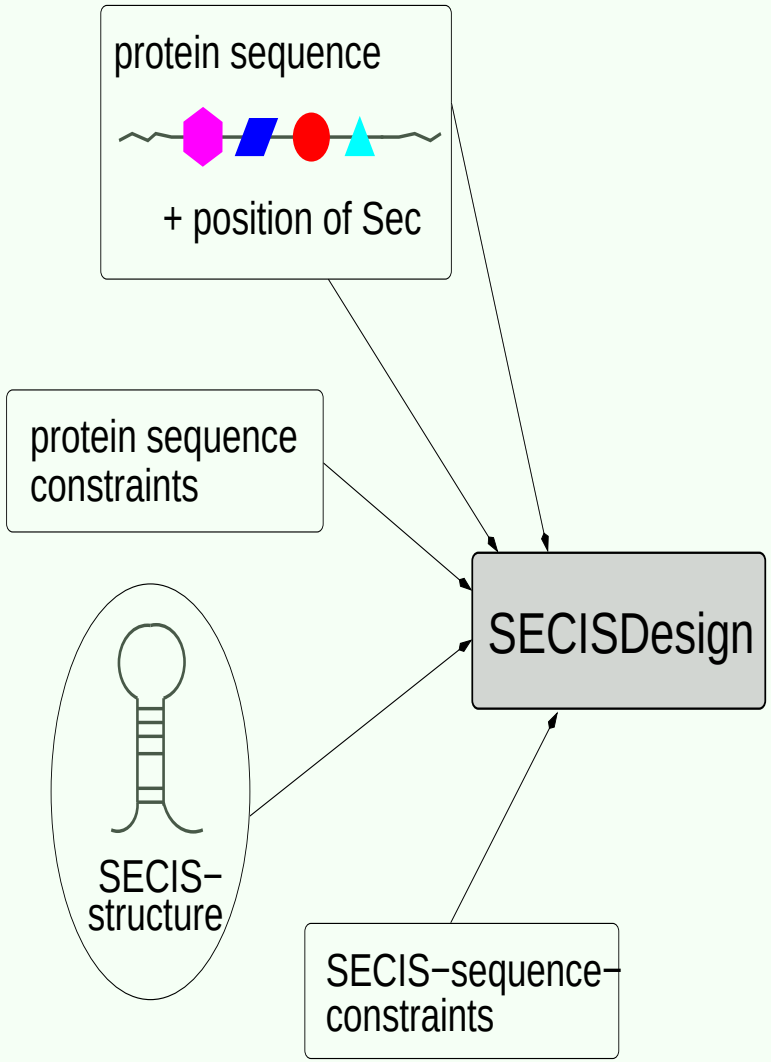
Extentions - 2

(2) Kinds of bonds in the SECIS-Element (mRNA)

- *bond* ... mandatory bond
- *optional bond* ... not necessary but of advantage, if formed
- *prohibited bond* ... a bond that is not allowed
- *unfavorable bond* ... not necessary but of advantage, if NOT formed



Input of SECISDesign



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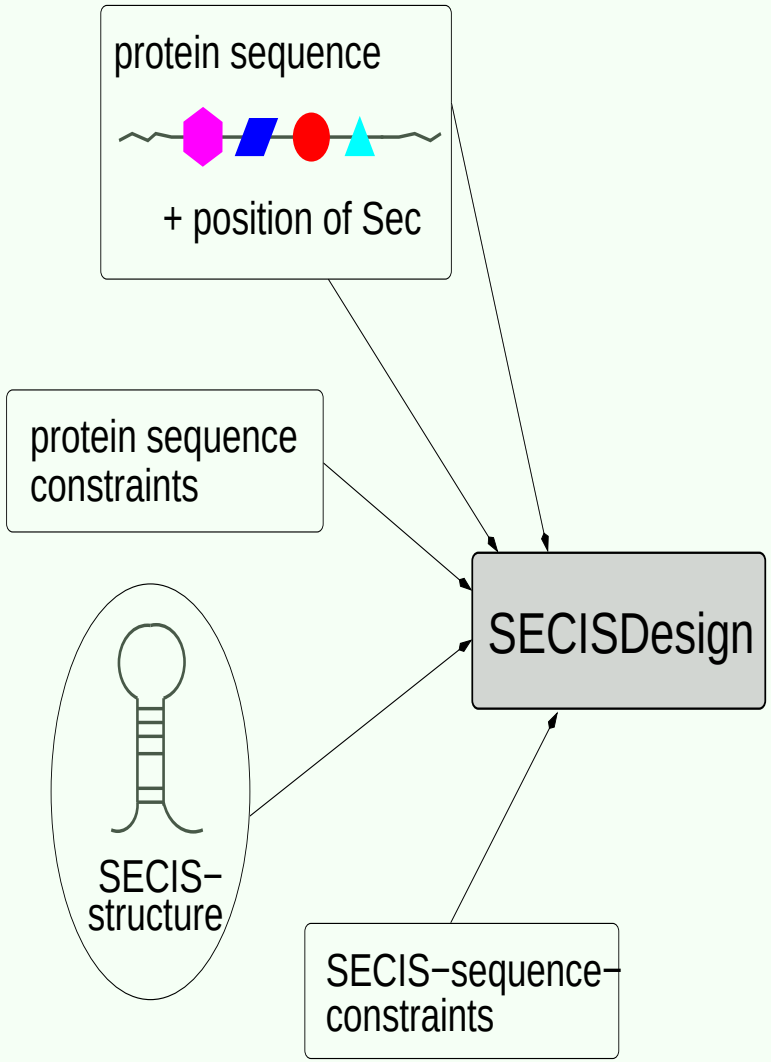

- **similarity functions** $f_i(L_i, a_i, t_i)$,

  where L_i ... *codon* corresponding to $N_{3i-2}N_{3i-1}N_{3i}$

  $t_i \in \{-1 \text{ (deletion)}, 0 \text{ (subst.)}, +1 \text{ (insertion)}\}$

  $a_i = \sum_{j=1}^{i} t_j$

  including:
  - *similarity at nucleotide level*
  - *similarity at amino acid level*

The Algorithm - Overview

Two Steps:

1. *Dynamic Programming:* (a) divides the problem in subproblems
   (b) solves subproblems and stores results
   (c) combines the sub-solutions

   - fi nds an *optimal* RNA sequence conc. *similarity* (nucleotide + amino acid) that can fold into the *SECIS-structure*
     (maximizes $\sum_{i=1}^{n} f_i(L_i, a_i, t_i)$)

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2. *Local Search:*

   - further improve the designed sequence to *increase the folding probability (inverse RNA folding)*

Dynamic Programming - 1

Central Function:

$$\begin{bmatrix} L_i & \ldots & \text{assignment of the } i\text{-th codon} \\ l & \ldots & \#\text{insertions - \#deletions left} \\ s & \ldots & \#\text{insertions - \#deletions inside interval} \end{bmatrix}$$

$$w_j^i(L_i, L_j, m, l, s) = \max_{\substack{L_{i+1} \ldots L_{j-1} \\ t(s)}} \left\{ \sum_{i<k\leq j} f_k(L_k, l + \sum_{g=1}^{k-i} t_g, t_{k-i}) \;\middle|\; \begin{array}{l} L_i \ldots L_j \text{ satisfy SECIS-graph,} \\ m \text{ real. opt. bonds at } L_i \ldots L_j \end{array} \right\}$$

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Value of Interest:

$$\max_{\substack{L_1, L_n, m \\ |l| \leq 1, |s| \leq n-1}} \{w_n^1(L_1, L_n, m, l, s) + f_1(L_1, l, l)\}$$

$\Rightarrow$ solved by dynamic programming

Dynamic Programming - 2

Recurrence Theorem: $\mathbf{w^{i}_{i+k}}(L_i, L_{i+k}, m, l, s) =$

$$\begin{cases} -\infty & \text{if } L_i \text{ and } L_{i+k} \text{ contradict the SECIS-constraints} \\ \max\limits_{\substack{L_p \\ \text{splits } (m_1,m_2) \text{ of } m \\ \text{splits } (s_1,s_2) \text{ of } s}} \left(\mathbf{w}^{i}_{p}(L_i, L_p, m_1, l, s_1) + \mathbf{w}^{p}{}_{i+k}(L_p, L_{i+k}, m_2, l+s_1, s_2) \right) & \text{otherwise} \end{cases}$$

where $p \in \{i+1, ..., j-1\}$ ... farthest codon having a bond with the i-th codon, if there is none, $p = i+1$

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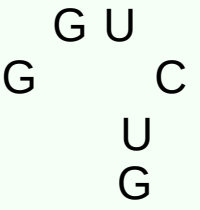

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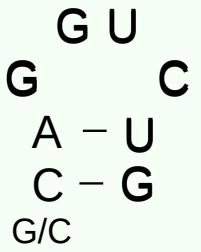

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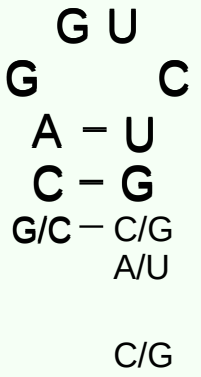

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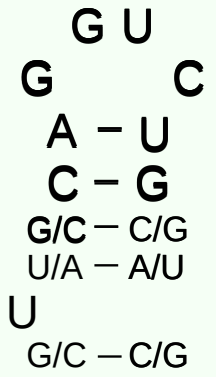

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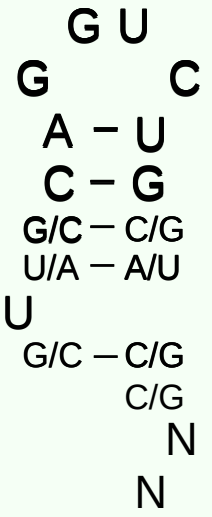

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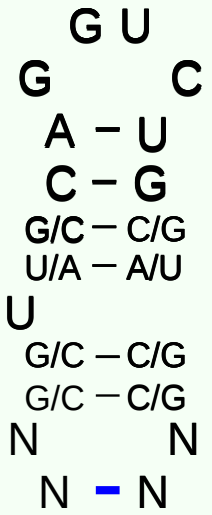

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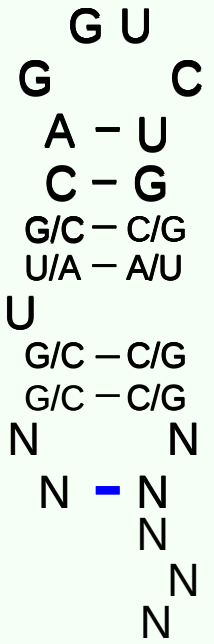

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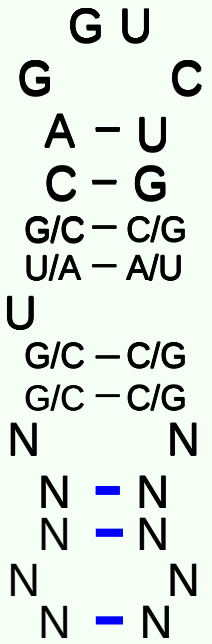

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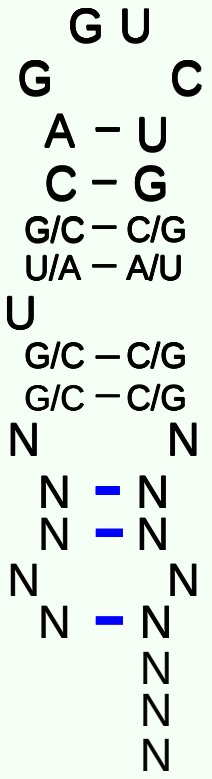

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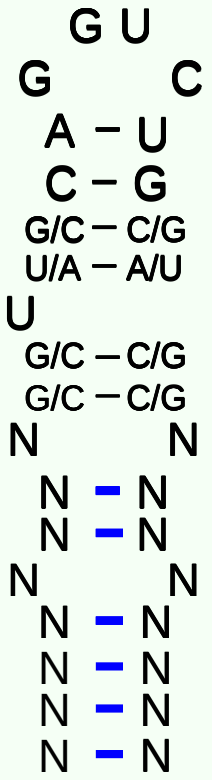

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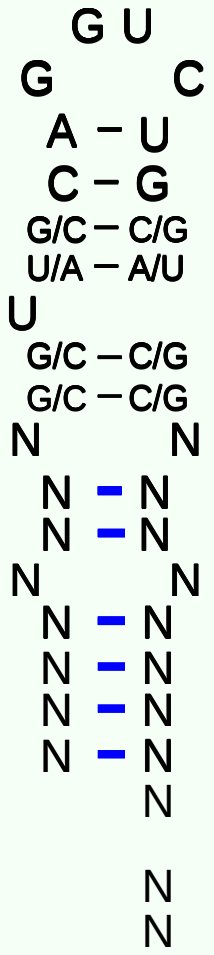

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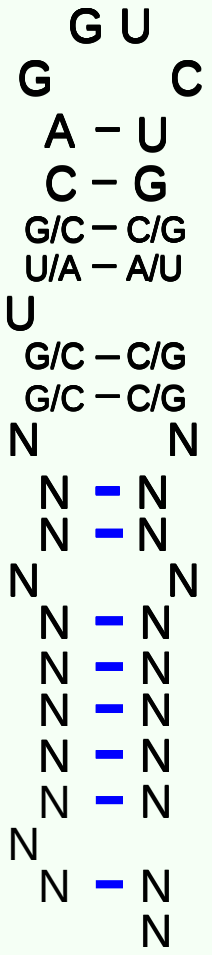

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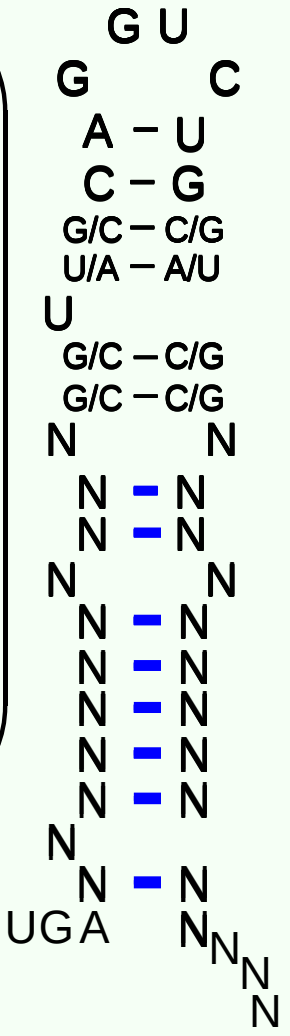

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After Step I:

Found a sequence that

- ***can*** fold into the SECIS-structure
- has ***maximal*** similarity to the SECIS-sequence
- encodes an amino acid sequence having ***maximal*** similarity to the original protein

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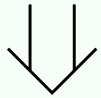


Step II: *(Local Search)*

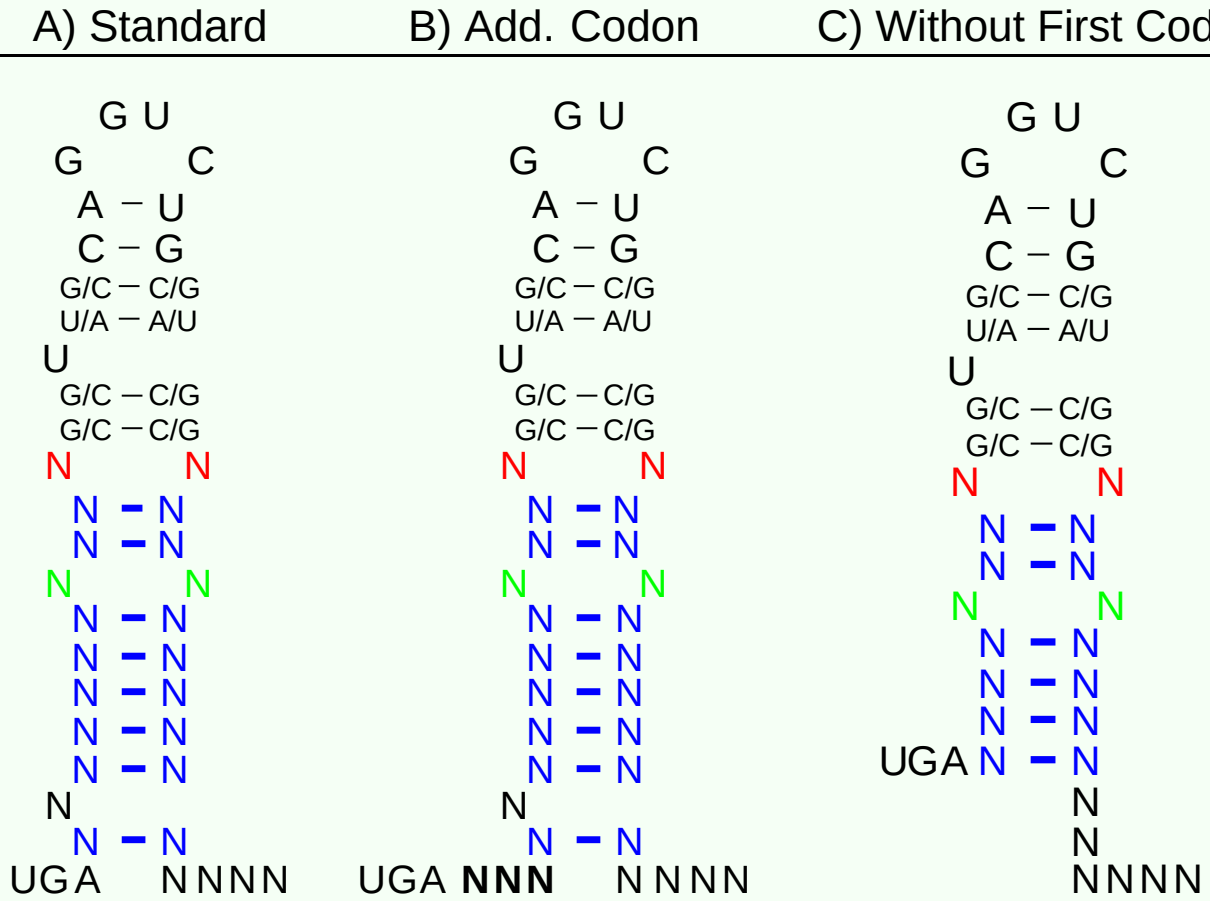
Local mutations to increase the folding probability (+ keep minimal similarities)

Methods:

- adaptive walk (as used in RNAinverse of the Vienna RNA Package)
- full local search
- stochastic local search

SECIS-elements

Alternatives of a SECIS-element in E.coli: (FdhF)



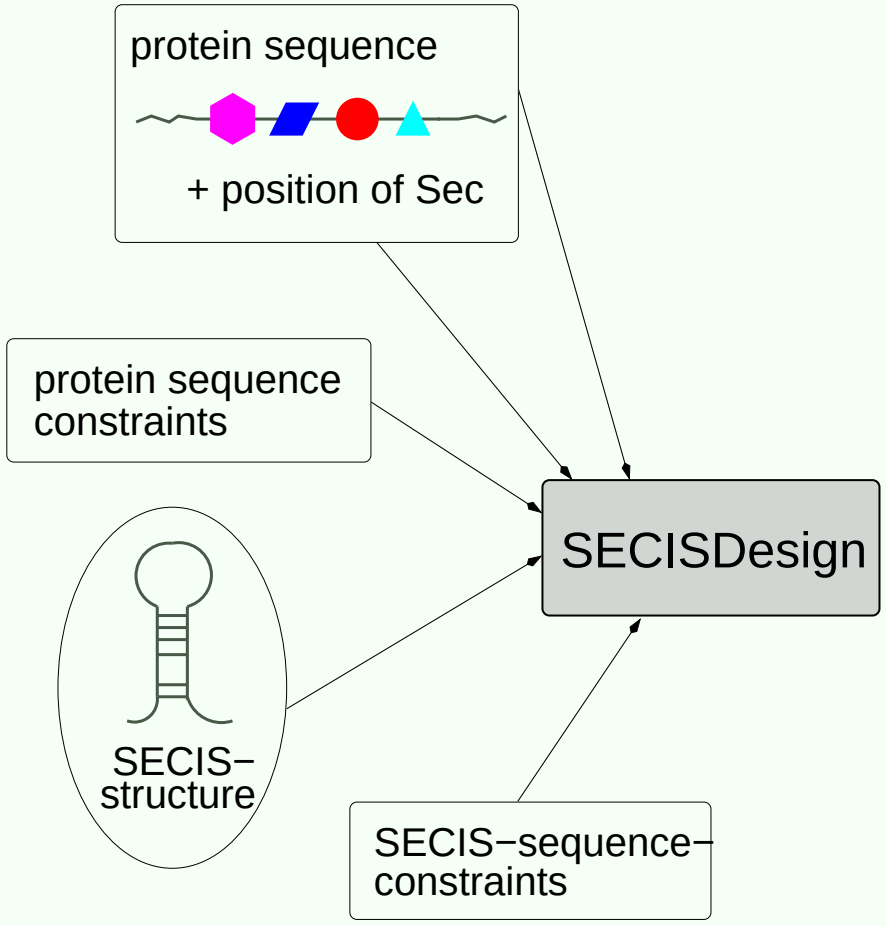
from Liu et al., NAR 1998

Results

Methionine Sulfoxide Reductase B (MsrB): (also analyzed in Bar-Noy et al., 2002)

	Amino Acid Sequences (starting after pos. of Sec)	*Sim. (BLOSUM62)*	*mRNA-Sequences and -Structures*	*Prob.*
mouse MsrB	**IFSSSLKFVPKGKE**			
MsrB_BaNo	**IFSTVAGLHPKGKE**	35	AUAUUUAGCACGGUUGCAGGUCUGCACCCUAAAGGCAAAGAA ....((.((((....))))))..............	0.01
mouse MsrB	**IFSSSLKFVPKGKE**			
MsrB_1A	**IFSSLPGLVPKGKE**	43	AUUUUCUCUUCGCUACCAGGUCUGGUGCCAAAAGGAAAAGAA .((((((.((.((.((((....)))))).)).))))))....	0.19
MsrB_2A	**IFSSLPGLVPQGAE**	33	AUCUUCUCGUCGCUACCAGGUCUGGUGCCACAAGGAGCCGAA ..((((((.((.((.((((....)))))).)).))))).....	**0.75**
mouse MsrB	**IFSSSLK-FVPKGKE**			
MsrB_1B	**IFSSSLPGLVPKGKE**	57+*IP*	AUAUUUUCCUCUUCGCUACCAGGUCUGGUGCCAAAAGGAAAAGAA ....((((((.((.((.((((....)))))).)).))))))....	0.08
MsrB_2B	**IVSSSLPGLVPQGAE**	40+*IP*	AUAGUCUCCUCGUCGCUACCAGGUCUGGUGCCACAAGGAGCAGAA ...(.((((((.((.((.((((....)))))).)).))))))....	**0.62**
mouse MsrB	**IFSSSLKFVPKGKE**			
MsrB_1C	**IFSLP-GLVPKGKE**	41+*DP*	AUCUUUUCGCUACCAGGUCUGGUGCCAAAAGGUAAAGAA (((((((.((.((((....)))))).)))))))....	**0.64**

The Algorithm - Summary



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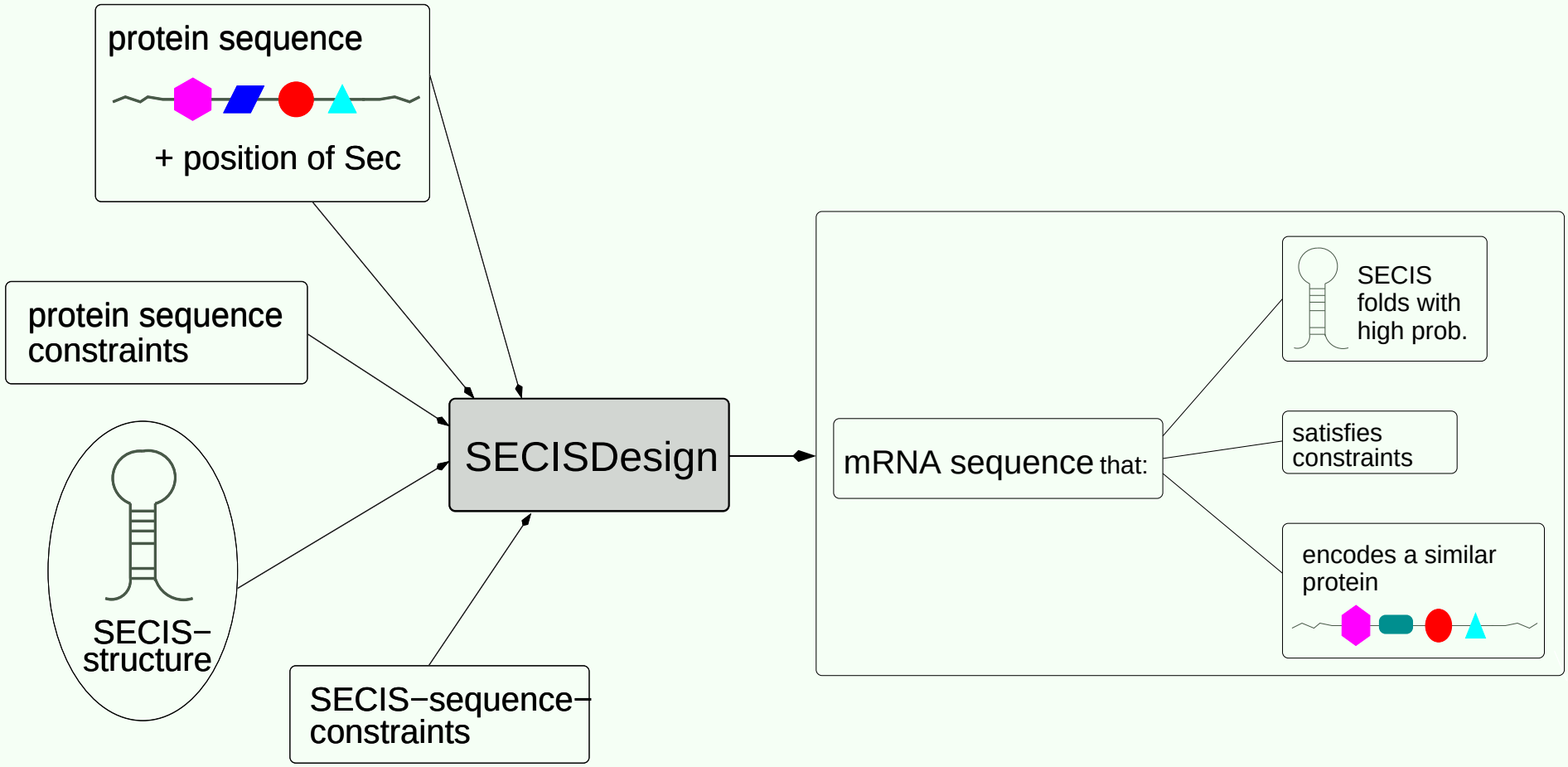
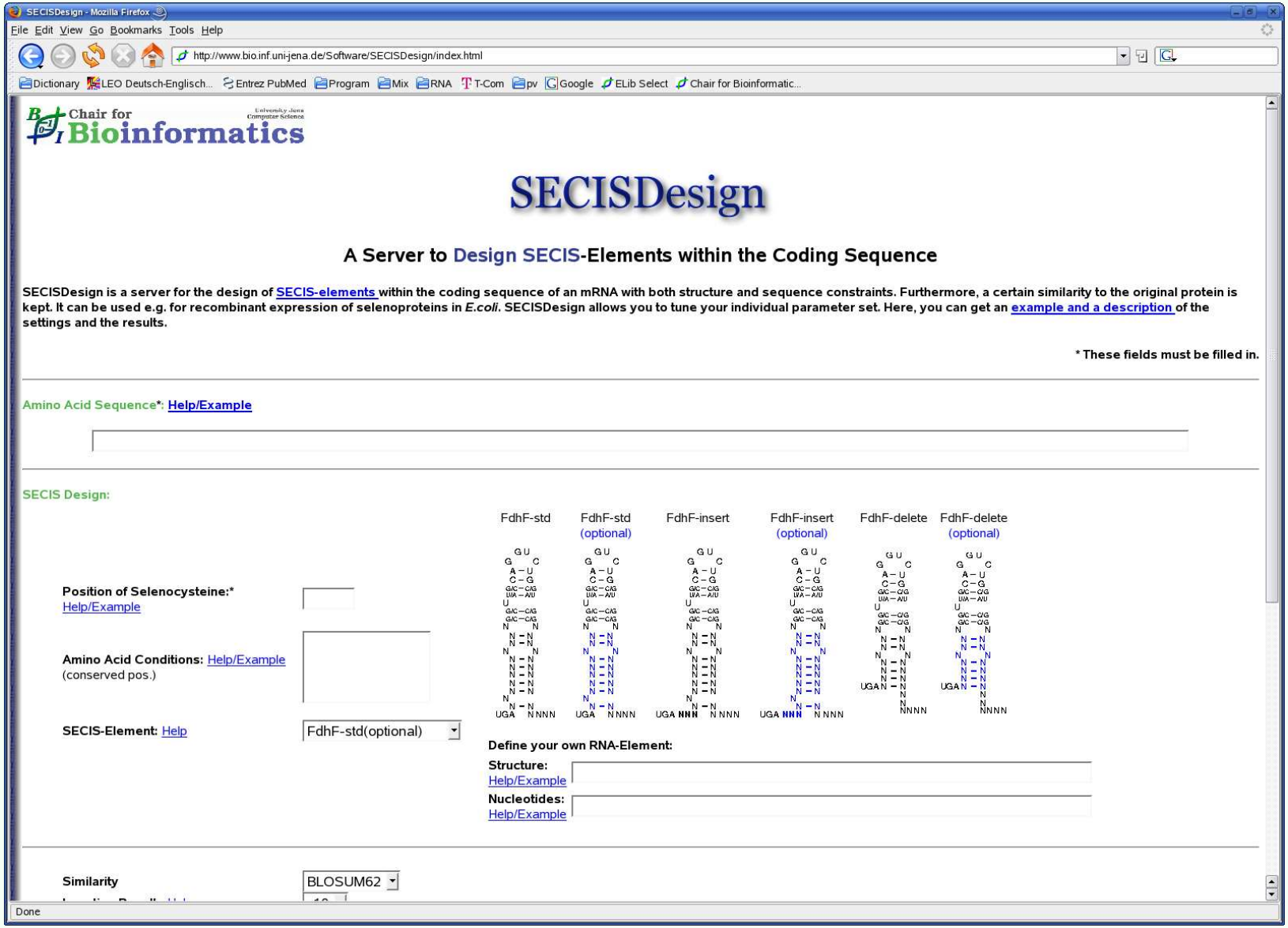


Webinterface

http://www.bioinf.uni-freiburg.de → Software → SECISDesign



Thank you for your attention.

Further reading:

Busch,A., S. Will, R. Backofen (2005). SECISDesign - A Server to Design SECIS-Elements within the Coding Sequence. *Bioinformatics*, **21(15)**, 3312-3.

Backofen, R. and Busch, A. (2004). Computational Design of New and Recombinant Selenoproteins. *Proc. of the 15th Annual Symposium on Combinatorial Pattern Matching (CPM2004)*.