# Who Cares? Graph Products in Theoretical Biology

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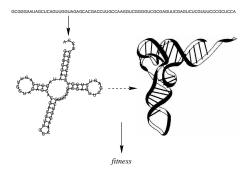
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#### Product Spaces and Biological Characters

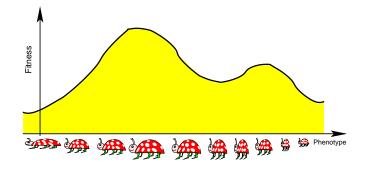
# **Evolutionary Biology: Genotype-Phenotype Maps**



#### Given:

- A set X of genotypes (sequences)
- Genetic Operators (mutation and/or recombination ...)
- A set Y of (potential) phenotypes (structures)
- A function f : X → Y assigning a phenotype to each genotype.

#### The "usual" view

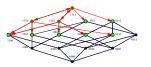


- Phenotypes are "somehow" numbers or vectors
- Accessible phenotypes are within a small (Euclidean) distance (in this vector space)
- Fitness is a (more or less) smooth function

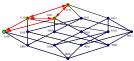
Population Genetics is perfectly happy ...

# So, what is wrong with this picture?

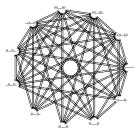
# (1) Genotype space is discrete



(2) Its structure depends on the genetic operator



(3) Phenotype space inherits its structure from genotype space



Accessibility at genotypic levels **implies** accessibility at phenotypic level **Does Evolution really "live" on an Euclidean space????** 

# Goal: A "Relative" Theory

We want a theory of phenotypes that can deal with concepts such as

- Continuity and Discontinuity
- Character
- Homology
- Innovation

#### WITHOUT recourse to a

specific representation of the phenotype

# **Genotype Spaces**

Given:

a set X of possible genotypes

a set A of realized genotypes

a fixed collection of genetic operators

[such as mutation, recombination, gene-rearrangement] **define** the set A' of genotypes accessible from A. Properties

- (i) No spontaneous creation, i.e,  $\emptyset' = \emptyset$ .
- (ii) A more diverse population produces more diverse offsprings:  $A \subseteq B$  implies  $A' \subseteq B'$
- (iii) All parental genotypes are also accessible in the next time step  $A \subseteq A'$ .

This is the same as the *C*-operator for the chemical network! In the case of mutation as the only source of diversity:

haploid populations, no sex, no recombination, etc

(iv) Diversity of offsprings depends only on the parent:

#### **Generalized Closure Spaces**

... instead of vector spaces ... Set X, closure function cl :  $\mathcal{P}(X) \rightarrow \mathcal{P}(X)$ Equivalent formulations:

$$\operatorname{int}(A) = X \setminus \operatorname{cl}(X \setminus A)$$

The interior is the dual of the closure function A set *N* is a neighborhood of *x* if and only if  $x \in int(N)$ . Let  $\mathcal{N}(x)$  be the set of all neighborhood of *x*.  $\mathcal{N} : X \to \mathcal{P}(\mathcal{P}(X))$ .

Closure, interior and neighborhood functions are equivalent.

#### **Generalized Closure Spaces**

	closure	neighborhood
K0	$cl(\emptyset) = \emptyset$	$X \in \mathcal{N}(x)$
K1	$A \subseteq B \implies \operatorname{cl}(A) \subseteq \operatorname{cl}(B)$	$N \in \mathcal{N}(\mathbf{x}), N \subseteq N'$
	$cl(A \cap B) \subseteq cl(A) \cap cl(B)$	$\implies$
	$cl(A) \cup cl(B) \subseteq cl(A \cup B)$	$N' \in \mathcal{N}(x)$
K2	$A \subseteq cl(A)$	$N \in \mathcal{N}(\mathbf{x}) \Rightarrow \mathbf{x} \in N$
K3	$cl(A\cup B)\subseteqcl(A)\cupcl(B)$	$N', N'' \in \mathcal{N}(\mathbf{x}) \implies$
		$N' \cap N'' \in \mathcal{N}(x)$
K4	cl(cl(A)) = cl(A)	$N \in \mathcal{N}(x) \iff$
		$int(N) \in \mathcal{N}(x)$
K5		$\mathcal{N}(\mathbf{x}) = \emptyset$ or
	$\bigcup_{i\in I} cl(A_i) = cl\left(\bigcup_{i\in I} A_i\right)$	$\exists N(x) : N(x) \subseteq N$
	iel (iel )	iff $N \in \mathcal{N}(x)$

In general: only (K0), (K1), (K2) hold: **neighborhood space** For mutation in haploid populations: (K0), (K1), (K2), (K5) [and thus (K3)]: **additive pretopological space** 

For comparison: (K0), (K1), (K2), (K3), and (K4) are equivalent to the axioms of a topology.

#### Cool!

... so, real evolution, genetic algorithms, evolution strategies, multi-objective optimization heuristics, genetic programming, etc., etc., live on a neighborhood space.

... for mutation only, it is even a pretopology. Thus:

Directed graphs and finite pretopological spaces are the same thing

Should we care that our closure function is NOT idempotent? NO, Eduard Čech in the 1960s wrote a big, fat textbook on point set topology, where he showed that pretty much everything works in pretopologies — thus you can do topology without every talking about open or closed sets. (Just the proofs get a bit more tedious without this convenience.)

### **Product Spaces**

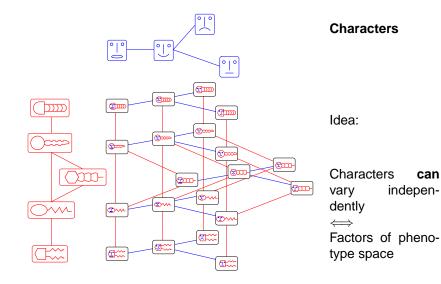
Let  $(X_1, c_1)$  and  $(X_2, c_2)$  be two general closure spaces, with neighborhood systems  $\mathcal{N}_1$  on  $X_1$  and  $\mathcal{N}_2$  on  $X_2$ . Product space: Point set  $X_1 \times X_2$ .

The neighborhoods of the product space satisfy:

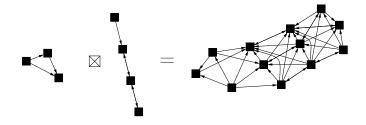
*N* is a neighborbood of  $(x_1, x_2)$  iff there are neighborhoods  $N_1 \in \mathcal{N}_1(x_1)$  and  $N_2 \in \mathcal{N}_2(x_2)$  such that  $N_1 \times N_2 \subseteq N$ . The projections  $\pi_i : (X_1 \times X_2, \mathcal{N}) \to (X_i, \mathcal{N}_i) : (x_1, x_2) \mapsto x_i$  are continuous functions for i = 1, 2. (as in topological spaces)

# What is a Phenotypic Character (Merkmal)?





#### 



Unique prime factor decomposition of connected graphs and digraphs.

Allows identification of global characters.

### **The Question**

#### Can we actually compute such factorizations???

For real data, i.e., when the original graph is known only appproximately?

"Regionally", i.e., when we want a coordinate system only is some part of the space?

 $\implies$  Wilfried's Presentation