# The Christmas Tree Dilemma 

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## Introduction



This is a Tree..

## Definitions

This is a Tree..


## Definitions

This is a Tree..

..and this is a Forest.


## Definitions

This is a Forest..
..and this is a Forest, too.


## Definitions

This is a Forest..
..and this is a Forest, too.

..and this..

..and this, too.


## Introduction II

This are christmas trees...


## Definitions

This is a Christmas tree..


## Definitions

This is a Christmas tree..

..and this, too.


## Definitions

This is a Christmas tree..

..and this, too.


This is a Christmas forest..

## Definitions

This is a Christmas tree..

..and this, too.


This is a Christmas forest..
..and this, too.


## Introduction III

These are Christmas ornaments..


## Introduction III

These are Christmas ornaments..



## Introduction III

These are Christmas ornaments..

..and this, too.


And this..


## Introduction III

These are Christmas ornaments..


And this..


..and this, too.


## Basics

## Assume that

- christmas tree $=$ tree
- christmas forest $=$ forest
- christmas ornament $=$ label



## Our Christmas Tree



## The Neighbours' Tree



Higher, better, nicer, stronger..


## Tree Editing

## Definition

A mapping between two ordered trees (forests) $T_{1}$ and $T_{2}$ is a binary relation on pairs of vertices $(x, y)$ and $\left(x^{\prime}, y^{\prime}\right)$ with $x, x^{\prime} \in V\left(T_{1}\right), y, y^{\prime} \in V\left(T_{2}\right)$ such that the following conditions hold:

- one-to-one condition: $x=x^{\prime} \Leftrightarrow y=y^{\prime}$
- ancestor condition: $x$ ancestor of $x^{\prime} \Leftrightarrow y$ ancestor of $y^{\prime}$
- sibling condition: $x$ left sibling of $x^{\prime} \Leftrightarrow y$ left sibling of $y^{\prime}$



## Tree Editing

Operations

- relabeling
- deletion
- insertion



## Tree Editing

## Operations

- relabeling
- deletion
- insertion


Aim: Minimize number of operations!
$\Rightarrow$ can be done with a DP algorithm.

## Tree Editing



## Tree Editing



## Tree Alignment

## Definition

A mapping of labels on forests $F_{1}$ and $F_{2}$ based on $(\mathcal{A} \cup\{-\}) \times(\mathcal{A} \cup\{-\})$ and restrictions $\pi_{1}(G)$ and $\pi_{2}(G)$ by considering either the first or the second coordinate, with $G$ being the resulting alignment forest.

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## Tree Alignment

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A mapping of labels on forests $F_{1}$ and $F_{2}$ based on $(\mathcal{A} \cup\{-\}) \times(\mathcal{A} \cup\{-\})$ and restrictions $\pi_{1}(G)$ and $\pi_{2}(G)$ by considering either the first or the second coordinate, with $G$ being the resulting alignment forest.


The cost of the alignment is the sum of the cost of label pairs. $\Rightarrow$ can be done with a DP algorithm.

## Tree Alignment



## Tree Alignment



## Grammars

## Definition

Context-free grammars with production rules of the form $V \rightarrow \alpha$, where $V$ is a non-terminal and $\alpha$ is a string of terminals and/or non-terminals.

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grammar + scoring algebra + index structure $=$ DP over arbitrary data structures

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grammar + scoring algebra + index structure $=$ DP over arbitrary data structures

## The Task

(1) find a grammar describing your problem
(2) find a scoring algebra
(0) find an index structure and iteration rules for your data structure (if not done yet)

## Grammars

Traversing Forests


## Grammars

## Tree Editing

|  | [start] |
| :---: | :---: |
| $\left.\begin{array}{l}F \\ F\end{array}\right) \rightarrow\binom{\epsilon}{\epsilon}$ | [end] |
| $\binom{F}{F} \rightarrow\binom{F}{F} \circ$ |  |
| $\left.\begin{array}{c}T \\ T\end{array}\right) \rightarrow\binom{F}{F}\binom{x}{\times}$ |  |
| $\left.{ }_{T}^{T}\right) \rightarrow\binom{F}{F}\left(\begin{array}{l}\bar{x} \\ \times\end{array}\right.$ |  |
| $\stackrel{T}{T}) \rightarrow\binom{F}{F}\binom{x}{-}$ | delin] |

The index structure is based on postorder of trees.

## Grammars

## Tree Alignment

$$
\begin{align*}
& \binom{S}{S} \rightarrow\binom{F}{F} \quad[s t a r t] \\
& \binom{F}{F} \rightarrow\binom{\epsilon}{\epsilon} \quad \text { [end] } \\
& \binom{F}{F} \rightarrow\binom{T}{T} \circ\binom{F}{F}[\text { iter] } \\
& \binom{T}{T} \rightarrow\binom{n}{n}\binom{F_{F}^{F}}{F} \quad \text { [align] } \\
& \binom{T}{T} \rightarrow\binom{-}{n}\binom{F}{F} \quad[\text { indel }] \\
& \binom{T}{T} \rightarrow\binom{n}{-}\binom{F}{F} \quad[d e l i n] \tag{1}
\end{align*}
$$

The index structure is based on preorder of trees.

## Why?

- DP algorithms on trees:
- small parsimony problem
- phylogenetic targeting
- tree editing
- tree alignment (with affine gap costs)
- automatized DP on various data structures (in future)
- inside/outside: DP on probabilities



## Challenge

Convert this title:
Why tree alignment doesn't have to suck

## Challenge

Convert this title:

## Why tree alignment doesn't have to suck

into something:

- referring to trees or forests
- but no other plants
- funny
- but not too funny ;)

The best submission will receive a christmas cucumber!

## Acknowledgements



> Thanks to...
> - Peter Stadler
> - Christian Höner zu Siederdissen

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- Sophia for pictures,
- Rojin for planting trees,
- Marc for the story of 'Weihnachtsgurke'


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- Sophia for pictures,
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- Marc for the story of 'Weihnachtsgurke'

Thank you for your attention!

## Proof for christmas cucumber!

LYRA-FAHRRAD-WERKE, HERMANN KLAASSEN, PRENZLAU
Erftes, älteftes, größtes und leiftungsfähigftes Spezial-Haus für Fahrräder, Sportartikel, Uhren, Goldwaren, Waffen, Spielwaren


Ns. 5046. Chriftbaumfthmud

 Karton mit 12 Stǜk...


Nr. 5047. Chriftbaumfimuck




