# Conservation in Long Non-coding RNAs and other Updates 

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## Incongruent Evolution

CGUGGAAACCCACAG
. (( ( (....)))) ..

CGUGGAAACC-CACAG

- ( ( ( (....) - ) ) ) . .
. ( ( $-((\ldots))))$.
CGU-GAAACCUCACAG

CGUGAAACCUCACAG
. ( ( ( ( . . . ) ) ) ) . .
. (( ( ( . . . ) ) ) ) . .
CGUGGAAACCCACAG
CGUGAAACCUCACAG
. (( ( (....)))) ..
exact conservation of the structure

## Incongruent Evolution

Similar sequence, shifted structure
..((( (. (( ( . . . .)))))..)))) .... AAGGCUCUAUUAACUGGUAUCGGCUAUAG ** * ***** ***** **** * * *** AAUGAUCUAUGAACUGUUAUCUGAUUUAG ...((().((( (....))))..))))...
..-((( (.((()....))))..))))....
AA-GGCUCUAUUAACUGGUAUCGGCUAUAG
** * * * * ***
AAUGAUCUAUGAACUGUUAUCUGAUU-UAG
...((( (. (( ( . . . .)))) ..)))) -...

## Is there really incongruent evolution?

## Sequence and structure alignments mir-30a and mir-30b

\author{
consensus ..U.AG....UGUAAACAUCCU. .ACU. . .AGCUGU.A. . .CA. . . .U.GGCU . . .A-GU. GGAUGUUUGC. .C. GC. . . CU mmu-mir-30a AGUGAGCGACUGUAAACAUCCUCGACUGGAAGCUGUGAAGCCACAAAUGGGCUUUCA-GUCGGAUGUUUGCAGCUGCCUACU
 mmu-mir-30b -UUCAGUUCAUGUAAACAUCCUACACU--CAGCUGUCAU--CAUGCGUUGGCUGGGAUGU-GGAUGUUUACGUCAGCUGUCU
 <br> consensus $\ldots$...[[[.[[[[[[[[[[[[........[[[[[.....................]]]].........]]]]]]]]]]]].]]].]... <br> 

## Incongruent Alignments

Basic idea: consider two or more alignments of the same objects (strings) simultaneously:

- implicitly defines alignments between the different copies of the same objects that do not allow mismatches
- Insertions and deletions in these same-object alignments correspond to shifts between the incongruent alignments
- scoring function:
weighted scores of the consitutent alignments + scores for the "shifts"
... what exactly are "shifts"?


## Formalization: Bi-Alignments

- two distinct alignments $\mathbb{U}$ and $\mathbb{V}$ of the same objects $\mathbf{a}$ and $\mathbf{b}$
- an alignment $\mathbb{W}$ of the columns of $\mathbb{U}$ and $\mathbb{V}$
- score $=u(\mathbb{U})+v(\mathbb{V})+w(\mathbb{W})$
- Bi-alignment Problem: simultaneously optimize $\mathbb{U}, \mathbb{V}$, and $\mathbb{W}$.



## Formalization: Shifts

- Gap patterns $c, d \in\binom{0}{1}$ in $\mathbb{U}$ and $\mathbb{V}$, respectively
- Congruent columns: $c_{1}=d_{1}$ and $c_{2}=d_{2}$.
- Incongruence $\|c-d\|=\left|c_{1}-d_{1}\right|+\left|c_{2}-d_{2}\right| \in\{0,1,2\}$
- score $w(\mathbb{W})$ : proportional to the sum of the incongruences of the alignment columns.
- An alignment of alignments is again an alignment:
$\mathbb{A} \simeq(\mathbb{U}, \mathbb{V}, \mathbb{W})$

- Number of in/dels between the two copies of $\mathbf{a}$ and $\mathbf{b}$ :

$$
d\left(\mathbb{A}_{13}\right)=\sum_{i}\left|c_{1}(i)-d_{1}(i)\right| \quad d\left(\mathbb{A}_{24}\right)=\sum_{i}\left|c_{2}(i)-d_{2}(i)\right|
$$

- Columnwise scoring of $\mathbb{A}$ :
score of the projected alignments $\mathbb{U} \simeq \mathbb{A}_{12}$ and $\mathbb{V} \simeq \mathbb{A}_{34}$ plus the in/del-only scores $d\left(\mathbb{A}_{13}\right)$ and $d\left(\mathbb{A}_{24}\right)$.


## Scoring Shifts

$$
\begin{gathered}
A \rightarrow A\left(\begin{array}{l}
\bullet \\
\vdots \\
\bullet
\end{array}\right)\left|A\left(\begin{array}{l}
\bullet \\
\vdots \\
-
\end{array}\right)\right| A\left(\begin{array}{l}
\bullet \\
\vdots \\
\bullet
\end{array}\right)|\cdots| A\left(\begin{array}{l}
- \\
- \\
-
\end{array}\right)\left|A\left(\begin{array}{l}
\bullet \\
- \\
-
\end{array}\right)\right| \varepsilon . \\
\left.\begin{array}{|c|cccc} 
& \left(\begin{array}{l}
\bullet \\
(\bullet)
\end{array}\binom{\bullet}{-}\right. & \left(\begin{array}{l}
- \\
(\bullet)
\end{array}\right. & (--) \\
(-) & \Delta & \Delta & \Delta & 2 \Delta \\
(-) & \Delta & 2 \Delta & 0 & \Delta \\
(-) & 2 \Delta & \Delta & \Delta & - \\
-
\end{array}\right) \\
M(0)=0 \\
M(x)=\max _{c \in \mathcal{C}} M(x-c)+s(x, c)
\end{gathered}
$$

## A very preliminary scan survey

- small and medium-width Rfam seed alignments
( $\leq 10$ sequences, $\leq 120$ columns)
1181 Rfam families
- check if Rfam consensus-structure "oriented" Rfam alignment is significantly different from a mafft re-alignment of the sequences 709 candidate families
- 10137 pairs of RNA sequences yield 143 cases in 72 families with prediced shifts in a sequence-based shift alignment


## Affine Gap Costs in $\mathbb{U}$ and $\mathbb{V}$

- Gotoh's algorithm for each of $\mathbb{U}$ and $\mathbb{V}$ scoring depending of the gap pattern of the penultimate column
- insufficient here: the penultimate column could be double-gap, i.e., an in/del of $\mathbb{W}$.
- remedy: keep end gap pattern defined for the last column that is not a double-gap:


## Including secondary structure

$$
\begin{gathered}
A \rightarrow A c|A \bar{c} A c| \varepsilon \\
u\left(\mathbb{U}, \varphi_{\mathbb{U}}\right)+v\left(\mathbb{V}, \varphi_{\mathbb{V}}\right)+w(\mathbb{W}) \\
M(x, y)=\max \left\{\begin{array}{l}
\max _{c \in \mathcal{C}} M(x, y-c)+s(y, c) \\
\max _{\substack{(z, y) \in \mathcal{B}^{*} \\
(c, d) \in \mathcal{C}^{*}}} M(x, z-c)+M(z, y-d)+\tilde{s}(z, c ; y, d)
\end{array}\right.
\end{gathered}
$$

$\mathcal{B}^{*} \ldots$ allowed index combinations, enforce base pairs

$$
(c, d) \in \mathcal{C}^{\prime}:=\left\{\left(\begin{array}{c}
- \\
- \\
\vdots
\end{array}\right),\left(\begin{array}{l}
\bullet \\
\vdots \\
\vdots
\end{array}\right),\left(\begin{array}{c}
\overline{-} \\
\vdots \\
\bullet
\end{array}\right),\binom{\bullet}{\vdots}\right\}^{2}
$$

... Sankoff-style Bi-Alignments

## Sankoff-style Bi-Alignments

- Complexity?
$O\left(n^{8}\right)$ entries times $O\left(n^{4}\right)$ operations
- BUT: number of shifts is limited:

$$
k \Delta \leq \delta^{*}(\mathbf{a}, \mathbf{b}):=\max _{\mathbb{U}} u(\mathbb{U})+\max _{\mathbb{U}} v(\mathbb{U})-\max _{\mathbb{U}}[u(\mathbb{U})+v(\mathbb{U})]
$$

reduction to $O\left(n^{4} k^{4}\right)$ entries with $O\left(n^{2} k^{2}\right)$ operations, i.e., $O\left(n^{6}\right)$ like the Sankoff algorithm

- locarna approximation: only include $O(n)$ most frequent base pairs for each structure
- reduction to $O\left(n^{2}\right)$ space and time.
- Implementation: on the way

Matrices $M_{p, q}$ indexed by end gap patterns $p$ and $q$ for $\mathbb{U}$ and $\mathbb{V}$

$$
M_{(p, q)}(x, y)=\max \left\{\begin{array}{l}
\max _{\substack{p^{\prime} \neq 0 \\
q^{\prime} \neq 0}} M_{\left(p^{\prime}, q^{\prime}\right)}(x-q, y-q)+s\left(\binom{x}{y},\binom{p^{\prime}}{q^{\prime}},\binom{p}{q}\right) \\
\max _{p^{\prime} \neq 0} M_{\left(p^{\prime}, q\right)}(x-p, y)+s\left(\binom{x}{y},\binom{p^{\prime}}{q},\binom{p}{0}\right) \\
\max _{q^{\prime} \neq 0} M_{p, q^{\prime}}(x, y-q)+s\left(\binom{x}{y},\binom{p}{q^{\prime}},\binom{0}{q}\right)
\end{array}\right.
$$

## Collaborators

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