

Convex Excess and Inequalities for Partial Cubes

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Joint work with [Sergey Shpectorov](#)

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- Hence partial cubes are precisely the graphs that admit isometric embeddings into hypercubes.

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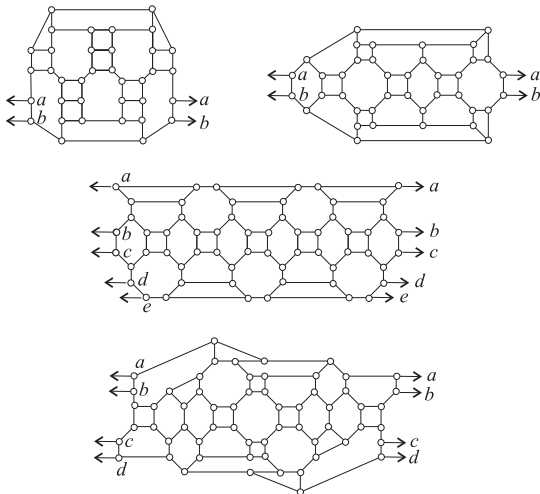
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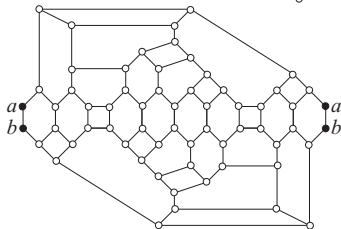
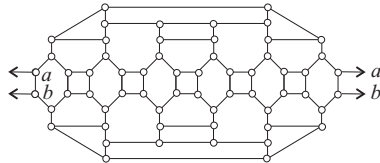
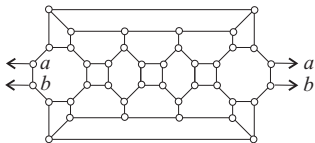
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- Cartesian products of partial cubes





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- G graph, new vertices its acyclic orientations, orientations differing by one edge of G .
- Integer partitions: vertices = partitions, edges = increment largest value and decrement some other value (or vice versa).
- Flips of triangulations; oriented matroids; media theory.

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Theorem (Djoković, 1973)

A connected graph G is a partial cube if and only if G is bipartite and for any edge uv of G the subgraph W_{ab} is convex.

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- Edges $e = xy$ and $f = uv$ of G are in **relation** Θ if
$$d(x, u) + d(y, v) \neq d(x, v) + d(y, u).$$

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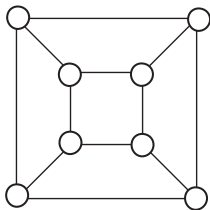
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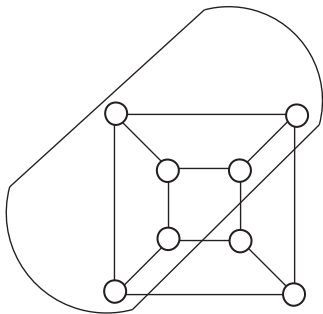
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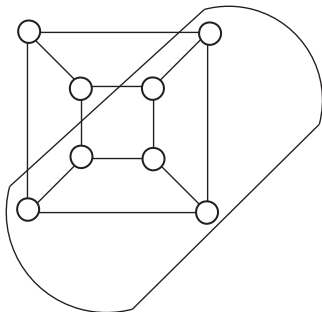
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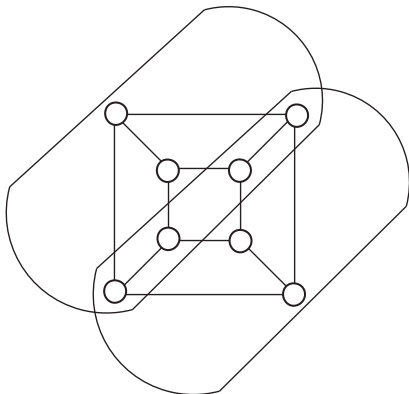
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- Insert the edges $v_1 u_1$ and $v_2 u_2$ whenever $v, u \in G_1 \cap G_2$ are adjacent in G .

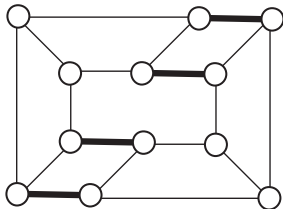
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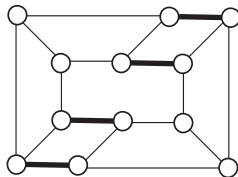
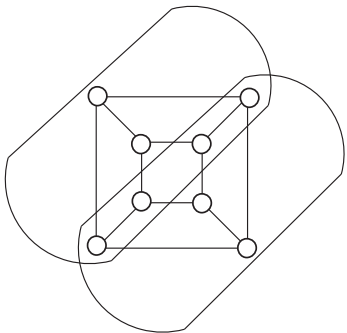
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That is, the smallest d such that G isometrically embeds into Q_d .

Other characterizations

Several other characterizations of partial cubes are known.

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- Bipartite ℓ_1 -graphs.
- Bipartite graphs whose distance matrix has exactly one positive eigenvalue.

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Two basic facts

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Median graphs and triangle-free graphs

Theorem (Imrich, K., Mulder, 1999)

Let $M(m, n)$ be the complexity of checking whether a graph G with m edges and n vertices is median. Then the complexity of checking whether G is triangle-free is at most $O(M(m, m))$.

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Theorem (Imrich, K., Mulder, 1999)

Let $T(m, n)$ be the complexity of finding all triangles of a given graph with m edges and n vertices. Then the complexity of checking whether a graph G on n vertices and m edges is a median graph is at most $O(m \log n + T(m \log n, n))$.

Inequality for median graphs

Theorem (K., Mulder, Škrekovski, 1998)

G median graph with n vertices and m edges. Then

$$2n - m - i(G) \leq 2.$$

Moreover equality holds if and only if G is Q_3 -free.

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Theorem (Brešar, K., Škrekovski, 2003)

Let G be a graph with n vertices and m edges that is obtained by a sequence of connected expansions from K_1 . Then $2n - m - i(G) \leq 2$. Moreover equality holds if and only if G is $C_t \square K_2$ -free ($t \geq 3$) and K_4 -free.

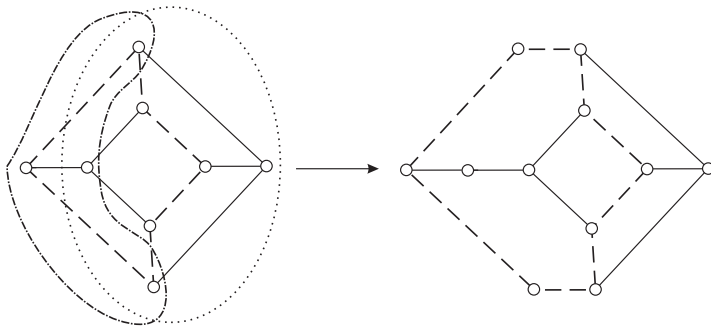
The question

Brešar, Imrich, K., Mulder, Škrekovski (JGT, 2002): Is there such an inequality for all partial cubes? In particular, does

$$2n - m - 2i(G) \leq 0$$

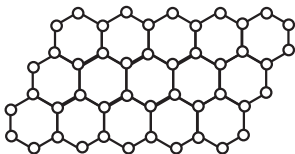
hold for any partial cube?

A reason for troubles



$2n - m - 2i(G) \leq 0$ need not hold

- $P(r, s)$, $1 \leq s \leq r$, parallelogram hexagonal graph.
- $n = (r + 1)(2s + 2) - 2$,
 $m = (r + 1)(2s + 1) - 2 + r(s + 1)$,
 $i(P(r, s)) = 2r + 2s - 1$.
- $2n - m - 2i(P(r, s)) = rs - 2(r + s) + 3$.



The inequality

- $\mathcal{C}(G) = \{C \mid C \text{ is a convex cycle of } G\}$. Then the **convex excess** of G :

$$ce(G) = \sum_{C \in \mathcal{C}(G)} \frac{|C| - 4}{2}.$$

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- **Spread partial cube**: all zone graphs are trees.

Theorem

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Moreover the equality holds if and only if G is a spread partial cube.

Proof

Proposition

Let G be a partial cube and let \tilde{G} be the expansion of G with respect to an isometric cover G_1, G_2 . If C is a convex cycle of G , then its expansion \tilde{C} is a convex cycle of \tilde{G} .

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Proposition

The zone graphs of partial cubes are connected.

Proof cont'd

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- $\tilde{n} = n + n_0$ and $\tilde{m} = m + n_0 + m_0$.

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- By the two propositions, \tilde{G} contains at least $t - 1$ convex cross cycles (with respect to G_1, G_2) of length at least six. So $ce(\tilde{G}) \geq ce(G) + t - 1$.

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- $m_0 \geq n_0 - t$.

Proof cont'd

$$\begin{aligned} & 2\tilde{n} - \tilde{m} - i(\tilde{G}) - ce(\tilde{G}) \\ & \leq 2(n + n_0) - (m + n_0 + m_0) - (i(G) + 1) - (ce(G) + t - 1) \\ & = (2n - m - i(G) - ce(G)) + (n_0 - m_0 - t) \\ & \leq 2 + (n_0 - (n_0 - t) - t) \\ & = 2. \end{aligned}$$

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The last two conditions imply that Z_F is a tree.

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Corollary

G spread partial cube, C and C' different convex cycles that are edges of Z_F . Then these cycles share no edges outside F .

Proof cont'd

- G spread partial cube, F Θ -class F , G_1 and G_2 connected components of $G \setminus F$.

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- By induction $2n_1 - m_1 - i(G_1) - ce(G_1) = 2$ and $2n_2 - m_2 - i(G_2) - ce(G_2) = 2$.

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- G_{10} is a forest (it is isomorphic to a subgraph of Z_F).

Proof cont'd

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$$\begin{aligned} \sum_{C \in E(Z_F)} \frac{|C| - 2}{2} &= \sum_{C \in E(Z_F)} (ce(C) + 1) \\ &= n_0 - 1 + \sum_{j=1}^{t-1} ce(C^{(j)}). \end{aligned}$$

Proof cont'd

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$$- \left(1 + i(G_1) + i(G_2) - (n_0 - 1) - \sum_{j=1}^{t-1} ce(C^{(j)}) \right)$$

$$- \left(ce(G_1) + ce(G_2) + \sum_{j=1}^{t-1} ce(C^{(j)}) \right)$$

$$= (2n_1 - m_1 - i(G_1) - ce(G_1)) + (2n_2 - m_2 - i(G_2) - ce(G_2)) - 2$$

$$= 2 + 2 - 2 = 2.$$