

FROM GENOMES TO SUPERGENOMES

HOW TO DEAL WITH BETWEENNESS FABIAN EXTERNBRINK SCADS & BIOINFORMATIK, LEIPZIG FABIAN@BIOINF.UNI-LEIPZIG.DE

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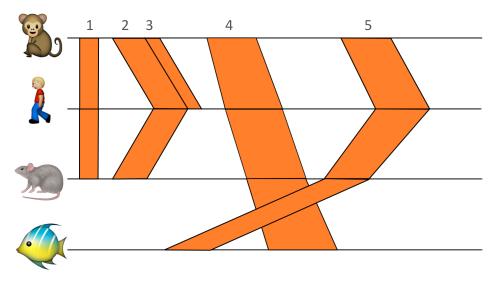
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GenomeRing: alignment visualization based on SuperGenome coordinates

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A Supergenome is a common coordinate system for all genomes in a multiple alignment.





- Multiple alignments
- Alignment blocks i.e. local best alignments
- Evolutionary events change the order
- Task:
 - Order the Blocks to create a common coordinate system



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TOTAL ORDERING PROBLEM*

J. OPATRNY[†]

Abstract. The problem of finding a total ordering of a finite set satisfying a given set of in-between restrictions is considered. It is shown that the problem is NP-complete.

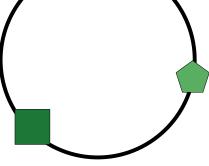
Key words. algorithms, computational complexity, total ordering, NP-completeness

Deciding Problem Given a finite set X and a collection $C \subseteq X^3$, is there a total order on X such that (i, j, k) \in C either i < j < k or i > j > k?

Optimization Problem

Given a finite set X and a collection $C \subseteq X^3$, find a maximal subset S from C, for which the decision problem w.r.t. S is true.

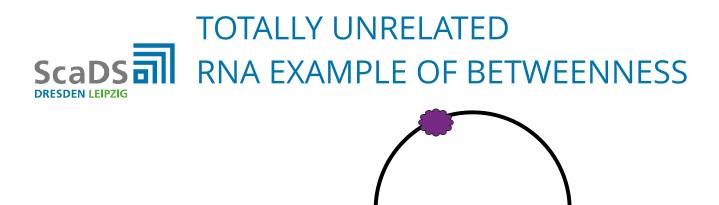


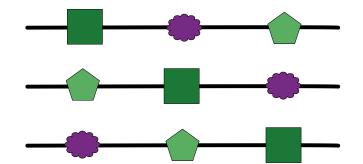


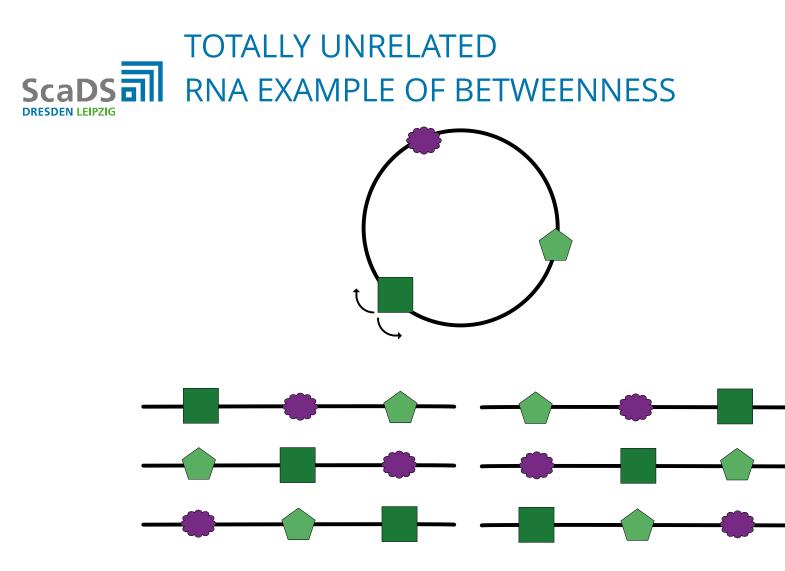
- Given a circular RNA molecule
- Different marker may exist on the RNA molecule
- Question:
 - What was the linear transcript?
 - Linear order of the marker?

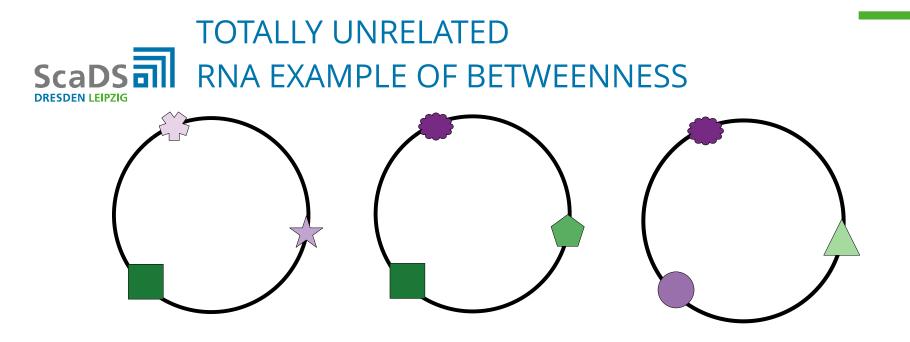




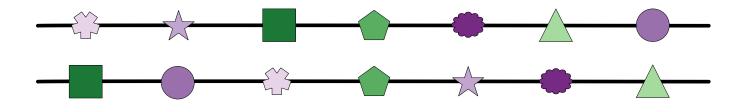


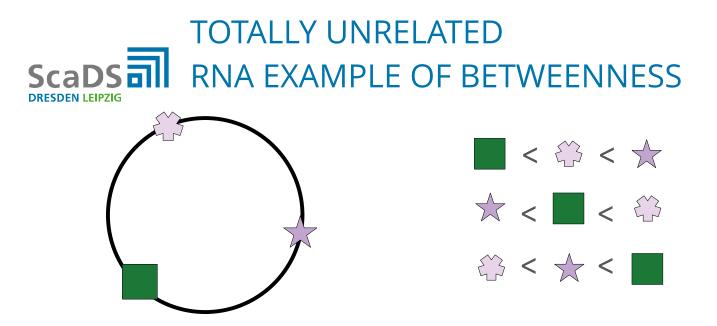






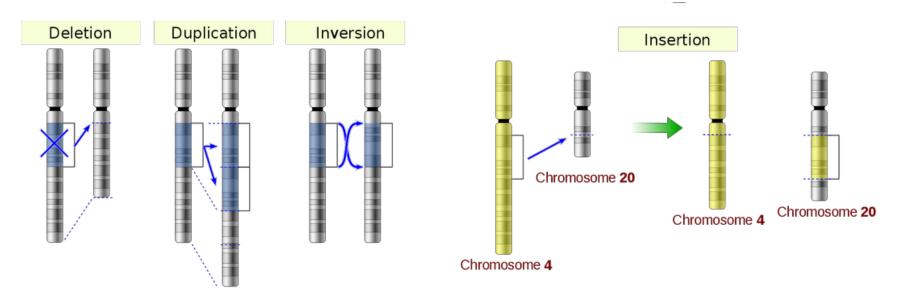
- More then one RNA molecule from one transcript
- Reading directions of the RNA molecules are independent
- Splicing allows deletions of markers





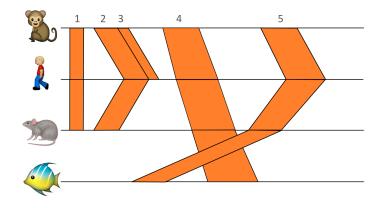
- It is a Betweenness Optimization Problem
 - The direction is not clear
 - Because of circularity, not all triples can be fulfilled
- Find largest subset of Triples that can be fulfilled by a linear order
- The linear order is the most likely linear transcript

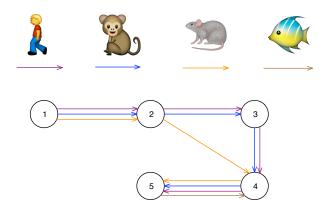




- Supergenome Problem is a Betweenness Optimization Problem
- Chromosome mutations
 - Direction of blocks is not clear
 - Not all triples can be fulfilled



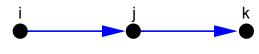




- Create graph from alignment
- One block is one vertex
- A edge from block v to w in color <u>x</u> is added if block w is successor of v in a genome <u>x</u>



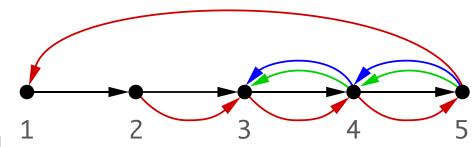
- Extends betweenness problem to a graph.
- Colored Multigraph Betweenness Problem
 - Find a maximal subset of colored edges E' of the multigraph such that the set of triples C(E') has a total order, where (i,j,k) ∈ C(E') if and only if there are two edges{i,j} and {j, k} with the same color.
- Idea to solve this is:

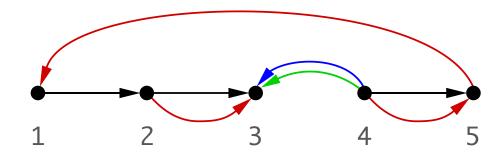


- Create a order of the vertices of the graph.
- Then calculate which edges are in the subset.



- Topological Sorting
 - Create order out of a DAG
 - If there is an edge from v
 to w, than v is before w in ¹
 the order
- Feedback Arc Sets (FAS)
 - Create a DAG
 - Remove as less edges as possible

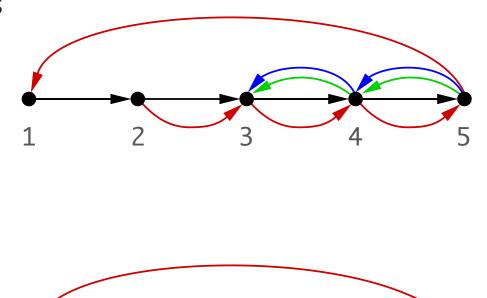




Order:4,5,1,2,3



- The resulting order destroys many betweenness information
- FAS does not fit betweenness problem well
 - Create artificial sinks and sources



(1,2,3) (4,5,1)



3

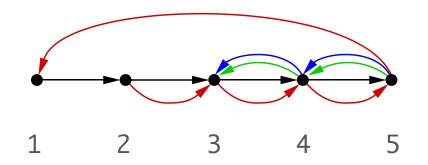
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1

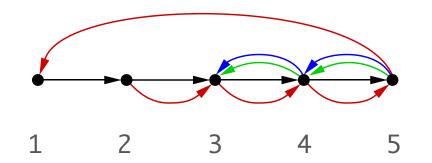
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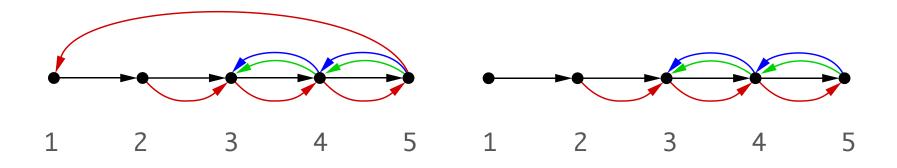
(1,2,3) (2,3,4) (3,4,5)(2,3,4) (3,4,5) (4,5,1)(5,4,3)(5,4,3)





$$(1,2,3)$$
 $(2,3,4)$ $(3,4,5)$
 $(2,3,4)$ $(3,4,5)$ $(4,5,1)$
 $(5,4,3)$
 $(5,4,3)$

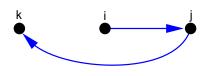




(1,2,3) (2,3,4) (3,4,5)(2,3,4) (3,4,5)(5,4,3)(5,4,3)



- Betweenness allow some cycles.
- Remove of all cycles is too much
- Two types of cycles:
 - Inconsistent cycles
 - Undirected cycles





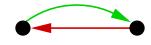
- If only inconsistent cycles in the graph
 - Solution to FAS and Betweenness Problem is the same



- Remove (most of) the undirected cycles
 - Mini-cycle Remover
- FAS is NP-complete
 - Use heuristic
- Noise reduction by simplifying collinear parts of the graph.
 - Sink/source simplifier
 - Closed-DAG simplifier

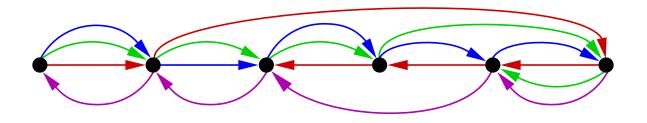


- Type of the cycle depends on used total order
- No order is given at this point
 - Use heuristic to find undirected cycles
- A mini-cycle with only two vertices
 - Very likely a undirected cycles
- Remove all mini-cycles in a intelligent way
 - If two mini-cycles share a vertex remove them together
 - Avoids generation of artificial sinks and sources

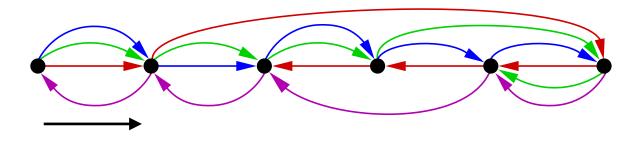




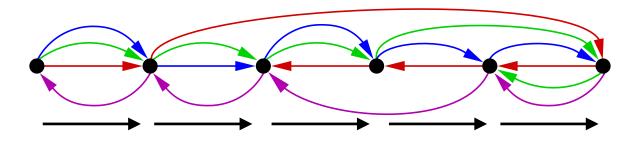




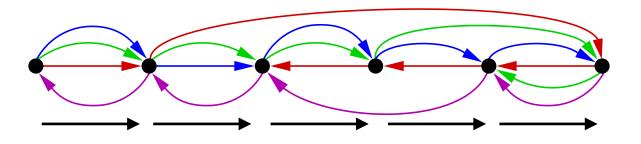


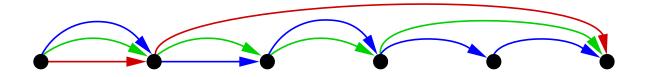






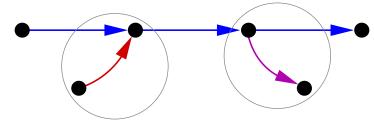








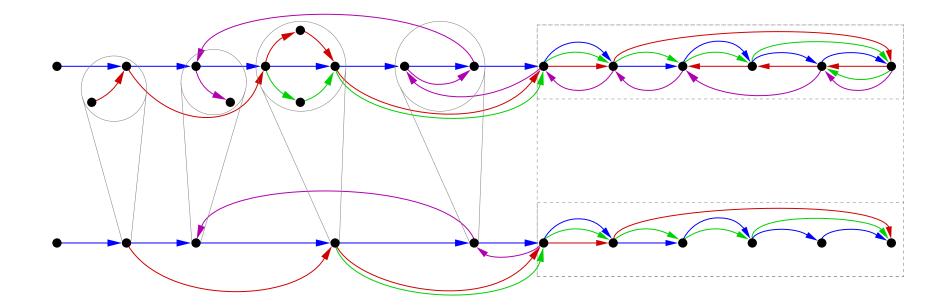
- A sink/source with only one predecessor/successor
- The position in the order is only influenced by the predecessor/successor
- It can be placed directly behind/before this predecessor/successor
- This is a collinear part in the graph
 - Simplified to one vertex





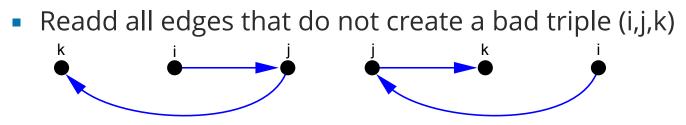
- A Closed-DAG is a collinear part in the graph
- It has this features:
 - It is a directed acyclic graph
 - It is connected to the rest of the graph by a single source vertex v and a single sink vertex w
 - All direct successors of v and all direct predecessors of w are contained in it
 - All vertices in it are successors of v and predecessors of w
- The Closed-DAG is an atomic unit in the order.







- Directed acyclic graph
 - Remove all edges that go from a vertex on position i to a vertex on position j if j<i/li>
- Betweenness graph
 - Add invers Edges



- Number of edges and triples can be counted
- No gold standard!
 - Can be compared with the start graph



- Two UCSC Datasets.
 - Created with a Reference Species
- Yeast
 - 7 species
 - 43495 vertices
 - 203275 edges, 197043 triples
- Insects
 - 27 species
 - 1451433 vertices
 - 25549792 edges, 25540919 triples



Edges % (triples %)	Yeast DAG	Insect DAG	Yeast Betweenness	Insect Betweenness
Simple FAS				
No mini-cycle Remover				
No simplifier				
All				



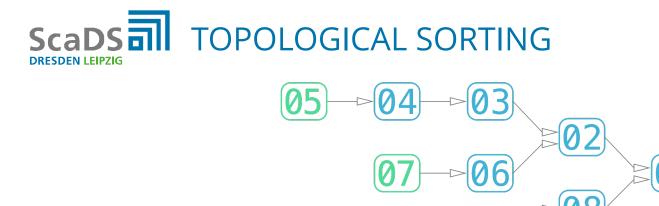
Edges % (triples %)	Yeast DAG	Insect DAG	Yeast Betweenness	Insect Betweenness
Simple FAS	66.87 (53.86)	61.54 (52.49)	82.65 (66.56)	86.97 (75.24)
No mini-cycle Remover	66.86 (53.87)	61.55 (52.50)	82.67 (66.60)	86.98 (75.26)
No simplifier	60.99 (59.29)	56.75 (56.37)	96.96 (94.04)	99.39 (98.80)
All	61.04 (59.31)	56.75 (56.37)	96.91 (93.94)	99.39 (98.80)



- Betweenness is everywhere!
- Solve optimization problem
- New graph based solution
 - Maximal subset of Edges
- Well studied approaches does not fit well
 - Can be fixed by a preprocessing
- Results can be measured
- Results look very promising



THANK YOU FOR YOUR ATTENTION



Topological sorting is not unambiguous

10

- Valid orders e.g.:
 - **5**,4,3,7,6,2,10,9,8,1
 - **5**,7,10,4,9,6,3,8,2,1
 - 10,7,5,6,4,3,2,9,8,1
 - **7.10,5,4,9,3,6,8,2,1**



- Use Distance information
- Next vertex in order is chosen by distance
- Not optimal for betweenness
- Valid orders e.g.:
 - **5**,4,3,7,6,2,10,9,8,1
 - 7,6,5,4,3,2,10,9,8,1
 - **1**0,9,8,76,5,4,3,2,1



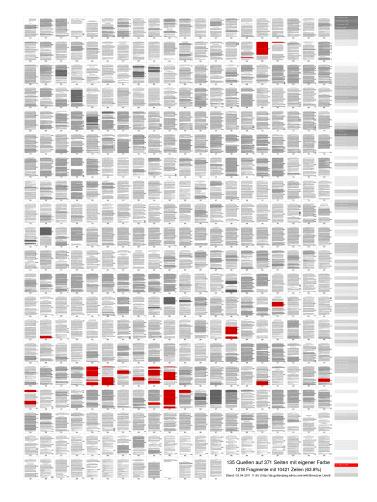
- Optimize minimal number of violation of the Robinson rule (1951): $\max(d(i,j),d(j,k)) \le d(i,k)$
- Change the order to an other valide topological sorting
- Check if number of violation is lowered
- Reaped until no further optimization is found
- **7,6**,5,4,3,2,10,9,8,1 → 5,4,3,**7,6**,2,10,9,8,1





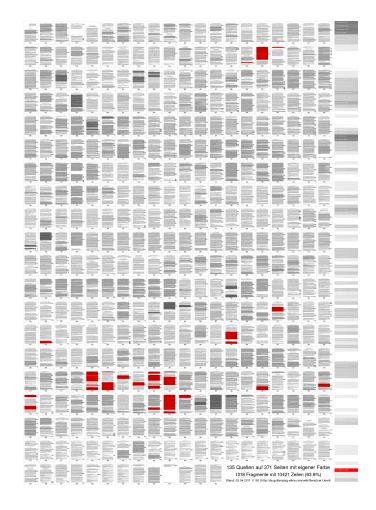
- 1218 plagiarism fragments
- 135 sources
- 63% of the work
- Sources widely distributed
- Possible questions:
 - Is basic structure from a source?
 - Which source is dominant in which part?

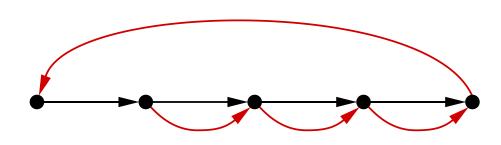




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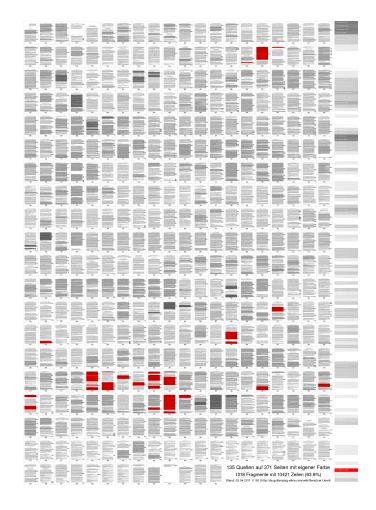


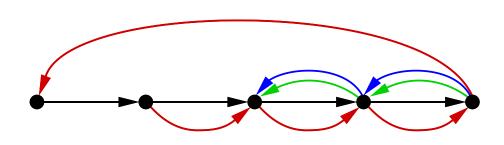




- Cites, pages, or sections as vertices
- Edges in the order of the dissertation and in order of the cites.



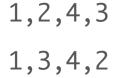


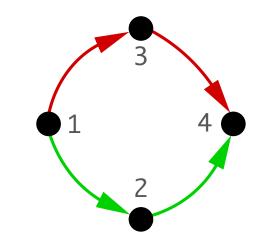


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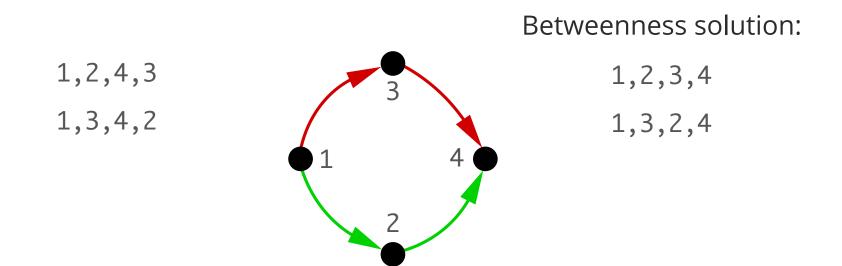
- A Hamiltonian path is a path that visits each vertex exactly once.
- The graph is connected
- Ignoring the direction of edges
 - Betweenness has no direction



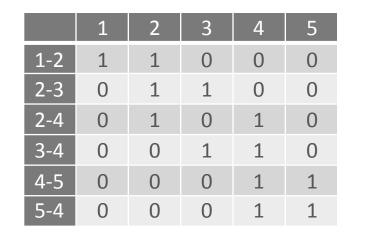


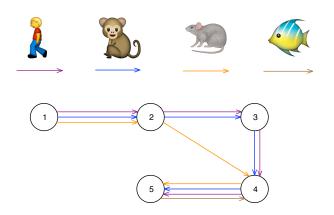


- Violated betweenness when two parts parallel
- Does not fit betweenness problem well



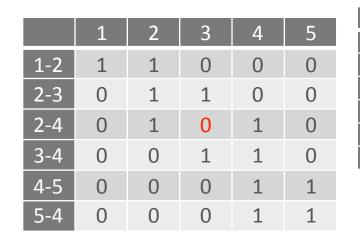


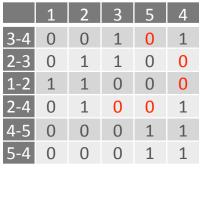


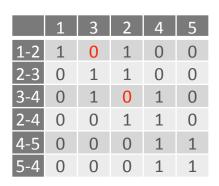


- Matrix with vertices as columns and adjacencies as rows
- Sort both the rows and columns of the matrix independently
- In such a way that rows and columns show all non-zero entries consecutively









- Consecutive ones property is violated even when betweenness is intact
- Bad adjacencies have huge impact
- Does not fit betweenness problem well