## ScaDS司| dRESDEN LEIPZIG

# FROM GENOMES TO SUPERGENOMES 

HOW TO DEAL WITH BETWEENNESS
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## ScaDS司 WHAT IS A SUPERGENOME? <br> DRESDEN LEIPZIG

## GenomeRing: alignment visualization based on SuperGenome coordinates

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A Supergenome is a common coordinate system for all genomes in a multiple alignment.

## ScaDS司 SUPERGENOME PROBLEM



- Multiple alignments
- Alignment blocks i.e. local best alignments
- Evolutionary events change the order
- Task:
- Order the Blocks to create a common coordinate system


## ScaDS司 WHAT IS BETWEENNESS? <br> dresden leipzig

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## TOTAL ORDERING PROBLEM*

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#### Abstract

The problem of finding a total ordering of a finite set satisfying a given set of in-between restrictions is considered. It is shown that the problem is $N P$-complete.


Key words. algorithms, computational complexity, total ordering, $N P$-completeness
Deciding Problem
Given a finite set $X$ and a collection $C \subseteq X^{3}$, is there a total order on X such that $(\mathrm{i}, \mathrm{j}, \mathrm{k}) \in \mathrm{C}$ either $\mathrm{i}<\mathrm{j}<\mathrm{k}$ or $\mathrm{i}>\mathrm{j}>\mathrm{k}$ ?

## Optimization Problem

Given a finite set $X$ and a collection $C \subseteq X^{3}$, find a maximal subset $S$ from $C$, for which the decision problem w.r.t. $S$ is true.

## ScaDS司 RNA EXAMPLE OF BETWEENNESS DRESDEN LEIPZIG



- Given a circular RNA molecule
- Different marker may exist on the RNA molecule
- Question:
- What was the linear transcript?
- Linear order of the marker?


## TOTALLY UNRELATED <br> ScaDS司 RNA EXAMPLE OF BETWEENNESS DRESDEN LEIPZIG



## TOTALLY UNRELATED <br> ScaDS司 RNA EXAMPLE OF BETWEENNESS DRESDEN LEIPZIG



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## ScaDS司 RNA EXAMPLE OF BETWEENNESS



- More then one RNA molecule from one transcript
- Reading directions of the RNA molecules are independent
- Splicing allows deletions of markers



# TOTALLY UNRELATED 

## ScaDS司 RNA EXAMPLE OF BETWEENNESS



- It is a Betweenness Optimization Problem
- The direction is not clear
- Because of circularity, not all triples can be fulfilled
- Find largest subset of Triples that can be fulfilled by a linear order
- The linear order is the most likely linear transcript


## ScaDS司 WHERE IS THE CONNECTION? <br> DRESDEN LEIPZIG



Insertion


Chromosome 20

Chromosome 4

- Supergenome Problem is a Betweenness Optimization Problem
- Chromosome mutations
- Direction of blocks is not clear
- Not all triples can be fulfilled


## ScaDS司 SUPERGENOME GRAPH




- Create graph from alignment
- One block is one vertex
- A edge from block $\mathbf{v}$ to $\mathbf{w}$ in color $\underline{x}$ is added if block w is successor of $\mathbf{v}$ in a genome $\underline{x}$


## ScaDS司 BETWEENNESS AND GRAPH? <br> DRESDEN LEIPZIG

- Extends betweenness problem to a graph.
- Colored Multigraph Betweenness Problem
- Find a maximal subset of colored edges E' of the multigraph such that the set of triples $C\left(E^{\prime}\right)$ has a total order, where $(i, j, k) \in C\left(E^{\prime}\right)$ if and only if there are two edges $\{i, j\}$ and $\{j, k\}$ with the same color.
- Idea to solve this is:

- Create a order of the vertices of the graph.
- Then calculate which edges are in the subset.


## ScaDS司 FAS AND TOPOLOGICAL SORTING <br> DRESDEN LEIPZIG

- Topological Sorting
- Create order out of a DAG
- If there is an edge from $\mathbf{v}$ to $\mathbf{w}$, than $\mathbf{v}$ is before $\mathbf{w}$ in 1
 the order
- Feedback Arc Sets (FAS)
- Create a DAG
- Remove as less edges as possible



## ScaDS司 FAS AND TOPOLOGICAL SORTING <br> DRESDEN LEIPZIG

- The resulting order destroys many betweenness information
- FAS does not fit betweenness problem well

- Create artificial sinks and sources
$(1,2,3) \quad(4,5,1)$



## ScaDS司 BETWEENNESS PROBLEMS



$$
\begin{aligned}
(1,2,3)(2,3,4) & (3,4,5) \\
(2,3,4) & (3,4,5)(4,5,1) \\
& (5,4,3) \\
& (5,4,3)
\end{aligned}
$$

## ScaDS司 BETWEENNESS PROBLEMS



$$
\begin{aligned}
&(1,2,3)(2,3,4) \\
&(2,4,5) \\
&(2,3,4)(3,4,5)(4,5,1) \\
&(5,4,3) \\
&(5,4,3)
\end{aligned}
$$

## ScaDS司 OPTIMAL BETWEENNESS SOLUTION



## ScaDS司 WHY IS FAS BAD?

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- Betweenness allow some cycles.
- Remove of all cycles is too much
- Two types of cycles:
- Inconsistent cycles

- Undirected cycles
- If only inconsistent cycles in the graph
- Solution to FAS and Betweenness Problem is the same


## ScaDS司 PREPROCCESING GRAPH <br> DRESDEN LEIPZIG

- Remove (most of) the undirected cycles
- Mini-cycle Remover
- FAS is NP-complete
- Use heuristic
- Noise reduction by simplifying collinear parts of the graph.
- Sink/source simplifier
- Closed-DAG simplifier


## ScaDS司 MINI-CYCLE REMOVER

- Type of the cycle depends on used total order
- No order is given at this point
- Use heuristic to find undirected cycles
- A mini-cycle with only two vertices

- Very likely a undirected cycles
- Remove all mini-cycles in a intelligent way
- If two mini-cycles share a vertex remove them together

- Avoids generation of artificial sinks and sources


## ScaDS司 MINI-CYCLE REMOVER EXAMPLE <br> DRESDEN LEIPZIG



## ScaDS司 MINI-CYCLE REMOVER EXAMPLE <br> DRESDEN LEIPZIG



## ScaDS司 MINI-CYCLE REMOVER EXAMPLE



## ScaDS司 MINI-CYCLE REMOVER EXAMPLE



## ScaDS司 SINK/SOURCE SIMPLIFIER

- A sink/source with only one predecessor/successor
- The position in the order is only influenced by the predecessor/successor
- It can be placed directly behind/before this predecessor/successor
- This is a collinear part in the graph
- Simplified to one vertex



## ScaDS司 CLOSED-DAG SIMPLIFIER

- A Closed-DAG is a collinear part in the graph
- It has this features:
- It is a directed acyclic graph
- It is connected to the rest of the graph by a single source vertex $\mathbf{v}$ and a single sink vertex $\mathbf{w}$
- All direct successors of $\mathbf{v}$ and all direct predecessors of w are contained in it
- All vertices in it are successors of $\mathbf{v}$ and predecessors of w
- The Closed-DAG is an atomic unit in the order.


## ScaDS司 THE BIG PICTURE DRESDEN LEIPZIG



## ScaDS司 MEASURE RESULTING ORDER <br> DRESDEN LEIPZIG

- Directed acyclic graph
- Remove all edges that go from a vertex on position i to a vertex on position j if $\mathrm{j}<\mathrm{i}$
- Betweenness graph
- Add invers Edges
- Readd all edges that do not create a bad triple (i,j,k)

- Number of edges and triples can be counted
- No gold standard!
- Can be compared with the start graph


## ScaDS司 DATASETS

- Two UCSC Datasets.
- Created with a Reference Species
- Yeast
- 7 species
- 43495 vertices
- 203275 edges, 197043 triples
- Insects
- 27 species
- 1451433 vertices
- 25549792 edges, 25540919 triples


## ScaDS司 RESULT <br> dresden leipzig

| Edges \% <br> (triples \%) | Yeast <br> DAG | Insect <br> DAG | Yeast <br> Betweenness | Insect <br> Betweenness |
| :--- | :--- | :--- | :--- | :--- |
| Simple FAS |  |  |  |  |

## ScaDS司 RESULT <br> DRESDEN LEIPZIG

| Edges \% <br> (triples \%) | Yeast <br> DAG | Insect <br> DAG | Yeast <br> Betweenness | Insect <br> Betweenness |
| :--- | :--- | :--- | :--- | :--- |
| Simple FAS | $66.87(53.86)$ | $61.54(52.49)$ | $82.65(66.56)$ | 86.97 (75.24) |
| No mini-cycle | 66.86 (53.87) | $61.55(52.50)$ | $82.67(66.60)$ | 86.98 (75.26) |
| Remover |  |  |  |  |

## ScaDS司 SUMMARY <br> DRESDEN LEIPZIG

- Betweenness is everywhere!
- Solve optimization problem
- New graph based solution
- Maximal subset of Edges
- Well studied approaches does not fit well
- Can be fixed by a preprocessing
- Results can be measured
- Results look very promising


## ScaDS司 <br> DRESDEN LEIPZIG

## THANK YOU FOR YOUR ATTENTION

## ScaDS司 TOPOLOGICAL SORTING



- Topological sorting is not unambiguous
- Valid orders e.g.:
- 5,4,3,7,6,2,10,9,8,1
- 5,7,10,4,9,6,3,8,2,1
- 10,7,5,6,4,3,2,9,8,1
- 7.10,5,4,9,3,6,8,2,1


## ScaDS司 DISTANCE TOPOLOGICAL SORTING



- Use Distance information
- Next vertex in order is chosen by distance
- Not optimal for betweenness
- Valid orders e.g.:
- 5,4,3,7,6,2,10,9,8,1
- 7,6,5,4,3,2,10,9,8,1
- 10,9,8,76,5,4,3,2,1


## ScaDS司 BETWEENNESS SORTING



- Optimize minimal number of violation of the Robinson rule (1951):

$$
\max (d(i, j), d(j, k)) \leq d(i, k)
$$

- Change the order to an other valide topological sorting
- Check if number of violation is lowered
- Reaped until no further optimization is found
- 7,6,5,4,3,2,10,9,8,1 $\longrightarrow 5,4,3,7,6,2,10,9,8,1$


## ScaDS司 GUTTENPLAG



- 1218 plagiarism fragments
- 135 sources
- 63\% of the work
- Sources widely distributed
- Possible questions:
- Is basic structure from a source?
- Which source is dominant in which part?


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## ScaDS司 GUTTENPLAG GRAPH <br> DRESDEN LEIPZIG



- Cites, pages, or sections as vertices
- Edges in the order of the dissertation and in order of the cites.


## ScaDS司 GUTTENPLAG GRAPH <br> DRESDEN LEIPZIG



- Cites, pages, or sections as vertices
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## ScaDS司 HAMILTONIAN PATH

- A Hamiltonian path is a path that visits each vertex exactly once.
- The graph is connected
- Ignoring the direction of edges
- Betweenness has no direction
$1,2,4,3$
$1,3,4,2$



## ScaDS司 HAMILTONIAN PATH <br> DRESDEN LEIPZIG

- Violated betweenness when two parts parallel
- Does not fit betweenness problem well

Betweenness solution:
$1,2,4,3$
$1,3,4,2$


1,2,3,4
$1,3,2,4$

## ScaDS司 SIMULTANEOUS CONSECUTIVE ONES

|  | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $1-2$ | 1 | 1 | 0 | 0 | 0 |
| $2-3$ | 0 | 1 | 1 | 0 | 0 |
| $2-4$ | 0 | 1 | 0 | 1 | 0 |
| $3-4$ | 0 | 0 | 1 | 1 | 0 |
| $4-5$ | 0 | 0 | 0 | 1 | 1 |
| $5-4$ | 0 | 0 | 0 | 1 | 1 |



- Matrix with vertices as columns and adjacencies as rows
- Sort both the rows and columns of the matrix independently
- In such a way that rows and columns show all non-zero entries consecutively


## ScaDS司 SIMULTANEOUS CONSECUTIVE ONES <br> DRESDEN LEIPZIG

|  | 1 | 2 | 3 | 4 | 5 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $1-2$ | 1 | 1 | 0 | 0 | 0 |
| $2-3$ | 0 | 1 | 1 | 0 | 0 |
| $2-4$ | 0 | 1 | 0 | 1 | 0 |
| $3-4$ | 0 | 0 | 1 | 1 | 0 |
| $4-5$ | 0 | 0 | 0 | 1 | 1 |
| $5-4$ | 0 | 0 | 0 | 1 | 1 |


|  | 1 | 2 | 3 | 5 | 4 |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $3-4$ | 0 | 0 | 1 | 0 | 1 |  |  |  |  |  |  |  |
| $2-3$ | 0 | 1 | 1 | 0 | 0 |  |  | 1 | 3 | 2 | 4 | 5 |
| $1-2$ | 1 | 1 | 0 | 0 | 0 |  | $1-2$ | 1 | 0 | 1 | 0 | 0 |
| $2-4$ | 0 | 1 | 0 | 0 | 1 | 1 | $2-3$ | 0 | 1 | 1 | 0 | 0 |
| $4-5$ | 0 | 0 | 0 | 1 | 1 |  | $3-4$ | 0 | 1 | 0 | 1 | 0 |
| $5-4$ | 0 | 0 | 0 | 1 | 1 |  | $2-4$ | 0 | 0 | 1 | 1 | 0 |

- Consecutive ones property is violated even when betweenness is intact
- Bad adjacencies have huge impact
- Does not fit betweenness problem well

