

# Sampling RNA Secondary Structures with Pseudoknots using Analytic Combinatorics

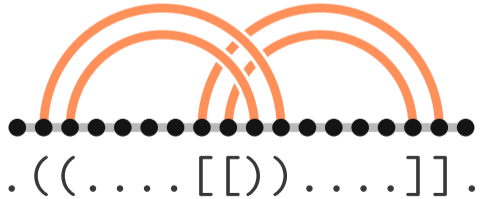
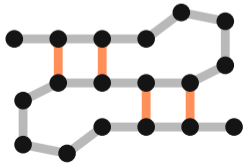
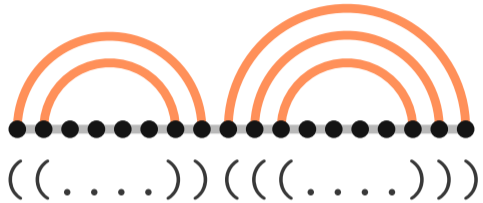
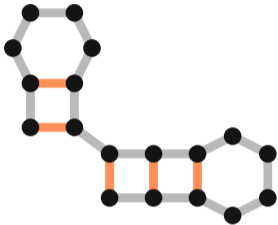
Casper Asbjørn Eriksen<sup>1</sup> Daniel Merkle<sup>1,2</sup>  
Markus Nebel<sup>2</sup> Jonas Vistrup<sup>1</sup>

<sup>1</sup>Algorithmic Cheminformatics Group, University of Southern Denmark

<sup>2</sup>Faculty of Technology, Bielefeld University

39th TBI Winterseminar

# RNA pseudoknots



# Grammars for RNA Secondary Structures

The following CFG generates well-formed dot-bracket strings (Motzkin words):

$$S \rightarrow \bullet S \mid (S) S \mid \epsilon$$

However, we might want a grammar with biologically meaningful rules.

- ◆ Energy estimations
- ◆ Parametrised generation
- ◆ *Uniform / non-uniform Boltzmann sampling*

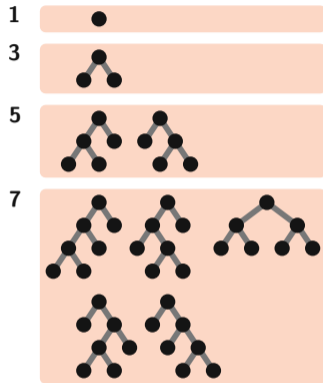
|   |   |
|---|---|
| $\left. \begin{array}{l} f_1 = S \rightarrow TAC, \\ f_2 = T \rightarrow TAC, \\ f_3 = T \rightarrow C, \\ f_4 = C \rightarrow C\bullet, \\ f_5 = C \rightarrow \epsilon, \\ f_6 = A \rightarrow (L), \\ f_7 = L \rightarrow (L), \\ f_8 = L \rightarrow M, \\ f_9 = L \rightarrow P, \\ f_{10} = L \rightarrow Q, \\ f_{11} = L \rightarrow R, \\ f_{12} = L \rightarrow F, \\ f_{13} = L \rightarrow G, \\ f_{14} = G \rightarrow (L)\bullet, \\ f_{15} = G \rightarrow (L)B\bullet\bullet, \\ f_{16} = G \rightarrow \bullet(L), \\ f_{17} = G \rightarrow \bullet\bullet B(L), \\ f_{18} = B \rightarrow B\bullet, \\ f_{19} = B \rightarrow \epsilon, \end{array} \right\} \begin{array}{l} \text{exterior loop} \\ \\ \text{initiate and extend stem} \\ \text{initiate multiloop} \\ \text{initiate interior loop} \\ \text{initiate hairpin loop} \\ \text{initiate bulge loop} \\ \\ \text{bulge loops} \end{array}$ | $\left. \begin{array}{l} f_{20} = F \rightarrow \dots, \\ f_{21} = F \rightarrow \dots, \\ f_{22} = F \rightarrow \dots\bullet\bullet\bullet H, \\ f_{23} = H \rightarrow H\bullet, \\ f_{24} = H \rightarrow \epsilon, \\ f_{25} = P \rightarrow \bullet(L)\bullet, \\ f_{26} = P \rightarrow \bullet(L)\bullet\bullet, \\ f_{27} = P \rightarrow \bullet\bullet(L)\bullet, \\ f_{28} = P \rightarrow \bullet\bullet(L)\bullet\bullet, \\ f_{29} = Q \rightarrow \bullet\bullet(L)K\bullet\bullet\bullet, \\ f_{30} = Q \rightarrow \dots J(L)K\bullet\bullet, \\ f_{31} = R \rightarrow \bullet(L)K\bullet\bullet\bullet, \\ f_{32} = R \rightarrow \dots J(L)\bullet, \\ f_{33} = J \rightarrow J\bullet, \\ f_{34} = J \rightarrow \epsilon, \\ f_{35} = K \rightarrow K\bullet, \\ f_{36} = K \rightarrow \epsilon, \\ f_{37} = M \rightarrow U(L)U(L)N, \\ f_{38} = N \rightarrow U(L)N, \\ f_{39} = N \rightarrow U, \\ f_{40} = U \rightarrow U\bullet, \\ f_{41} = U \rightarrow \epsilon. \end{array} \right\} \begin{array}{l} \text{hairpin loop} \\ \\ \text{small interior loops} \\ \\ \text{other interior loops} \\ \\ \text{multiloop} \end{array}$ |
|---|---|

(Nebel et. al 2011)

# Analytic Combinatorics

(Ordered) Binary trees:

$$B = \bullet + [\bullet \times B \times B]$$



*"If you can specify it,  
you can analyse it."  
- Philippe Flajolet*

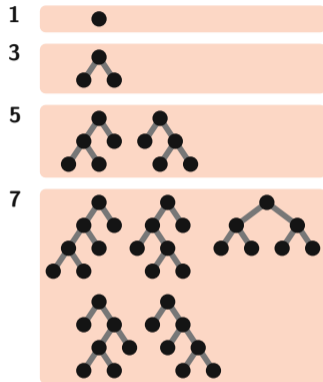
# Analytic Combinatorics

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Number of binary trees of size  $n$ :

1, 0, 1, 0, 2, 0, 5, 0, 14... (seq. **A000108**)



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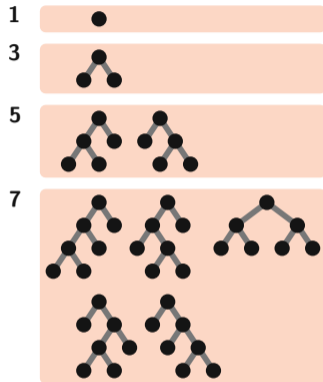
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**Generating function:**

$$GF_B(z) = z + zGF_B(z)^2 \quad \text{Symbolic transfer thm.}$$

$$= \frac{1 - \sqrt{1 - 4z}}{2} \quad \text{Solve + expansion}$$

$$= \mathbf{1}z + \mathbf{1}z^3 + \mathbf{2}z^5 + \mathbf{5}z^7 + \mathbf{14}z^9 + \dots$$

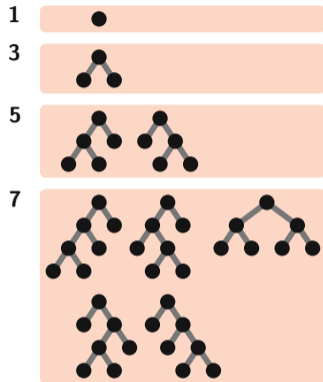


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# Analytic Combinatorics

Description of combinatorial structures:

- ◆ Binary trees:  $B = \bullet + [\bullet \times B \times B]$
- ◆ General trees:  $T = \bullet + [\bullet \times \text{SEQ}(T)]$
- ◆ Derangements:  $D = \text{SET}(\text{CYC}_{\geq 1}(\bullet))$



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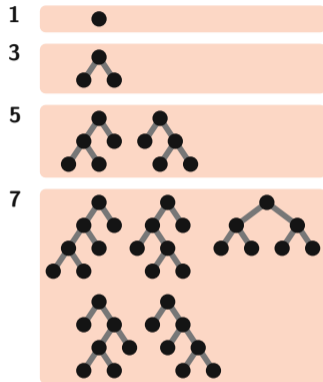
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A context-free language is a combinatorial class:

- ◆ Grammar:  $S \rightarrow \bullet S \mid (S) S \mid \epsilon$
- ◆ Specification:  $S = [\bullet \times S] + [(\times S \times) \times S]$



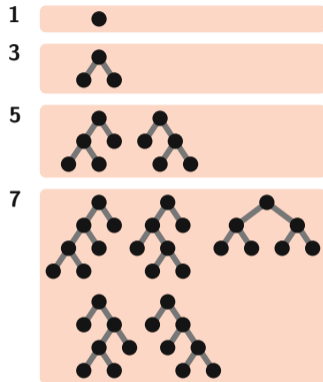
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# Analytic Combinatorics

A **Boltzmann Sampler** samples class  $\mathcal{A}$  uniformly

- ◆ Recursive algorithm using GF



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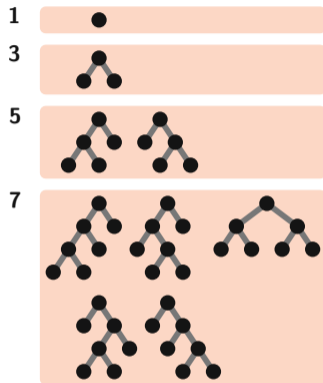
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## Maximum Likelihood Sampling

- ◆ Non-uniform sampling by weighing construction rules
- ◆ Obtain weights by parsing an ensemble of structures



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# Boltzmann Combinatorial Object Sampler

**BCOS** is a *work-in progress* Boltzmann sampling library for:

- ◆ Uniform sampling of arbitrary combinatorial classes
- ◆ Weighted sampling with predefined weights
- ◆ Maximum Likelihood ensemble-based sampling



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## Features:

- ◆ Dynamic arbitrary-precision evaluation of generating functions guarantees correctness
- ◆ Ability to sample complex classes even if no closed form GF can be found.
- ◆ Python interface for Combinatorial object semantics allows intuitive integration.

## Ordered tree

```
import bcos
tree_class = bcos.System(
    "T = N + (N * SEQ(T)) ,
    N = Atom"
)
tree_class.sample(size=(5,10))

> (N, ((N, (N, N, N, N)), (N, (N))))
```

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> (N, ((N, (N, N, N, N)), (N, (N))))
```

## RNA secondary structure

```
import bcos
g = bcos.cfg(
    "S -> . S | ( S ) | "
)
gclass = bcos.System(g)
gclass.sample(size=21)

> ..((.(..))..((.))(.))
```

# Multiple Context-Free Grammars

However, CFGs are not expressive enough to describe secondary structures with pseudoknots.

Introducing: **Multiple Context Free Grammars** (MCFG)

- ◆ Describes a broader class of languages than CFGs
- ◆ Allows for non-local correlations

Next step: designing an MCFG for pseudoknot structures.

Example MCFG  
( $a^n b^n c^n d^n$ ):

$$\begin{pmatrix} A_1 \\ A_2 \end{pmatrix} \rightarrow \begin{pmatrix} ab \\ cd \end{pmatrix}$$

$$\begin{pmatrix} A_1 \\ A_2 \end{pmatrix} \rightarrow \begin{pmatrix} aA_1b \\ cA_2d \end{pmatrix}$$

Example derivation:

$$\begin{aligned} & A_1 A_2 \\ \rightarrow & aA_1bcA_2d \\ \rightarrow & aabbccdd \end{aligned}$$

Problem: Potentially infinite alphabet

- ◆  $()$ ,  $[]$ ,  $\{\}$ ,  $\langle \rangle$ ,  $Aa$ ,  $Bb$ ,  $\dots$   $\dots$ ,  $A\alpha$ ,  $B\beta$ ,  $\Gamma\gamma$ ?



# Pseudoknot Grammars

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Possible solution:

- ◆ Parametrised alphabet:

$$S \rightarrow A_1A_2 \mid \bullet \mid ({}_X S)_X$$

$$\begin{pmatrix} A_1 \\ A_2 \end{pmatrix} \rightarrow \begin{pmatrix} SA_1 \\ A_2 \end{pmatrix} \mid \begin{pmatrix} A_1 \\ SA_2 \end{pmatrix} \mid \begin{pmatrix} ({}_X A_1) \\ ({}_X A_2) \end{pmatrix} \mid \begin{pmatrix} \epsilon \\ \epsilon \end{pmatrix}$$

where  $({}_X, )_X = ()$ ,  $[]$ ,  $\dots$ .

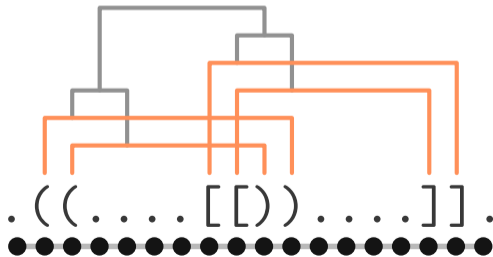
# Pseudoknot Grammars

Parsing:

- ◆ Instantiate grammar with sufficient degree and parser  
**or:**
- ◆ Generate the lowest level of the parse tree and parse this

Generation:

- ◆ Generate using simple grammar and create string from parse tree



## Strategy

- ◆ Let  $G$  be a parametrised MCFG (**PMCFG**)
- ◆ Let  $L$  be a library of pseudoknot RNA dot-bracket strings, possibly containing pseudoknots

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# Sampling Pseudoknot Structures

## Strategy

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- ◆ Result is a stochastic parametrised MCFG (**SPMCFG**)
- ◆ Convert to a weighted combinatorial class  $R$
- ◆ Calculate GF based on **SCFG** reduction
- ◆ Sample  $R$ , and create dot-bracket string based on the parse tree.



Thank you!

MATOMIC

novo nordisk  
**foundation**