

Bubble Space on Hypergraphs

Josef Leydold

WU Vienna · Austria

41st TBI Winterseminar
Bled – February 9, 2026

Motivation

Question by Michael Hecht:

(many years ago in a pub in Bled)

- ▶ Can we compute **higher Betti numbers** for graphs?

Motivation

Question by Michael Hecht:

(many years ago in a pub in Bled)

- ▶ Can we compute **higher Betti numbers** for graphs?
- ▶ What are these Betti numbers?
- ▶ For what are they good for?

Motivation

Question by Michael Hecht:

(many years ago in a pub in Bled)

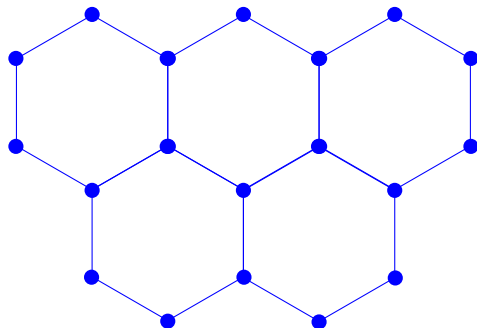
- ▶ Can we compute **higher Betti numbers** for graphs?
- ▶ What are these Betti numbers?
- ▶ For what are they good for?

Simple undirected **Graph** $G = (V, E)$

- ▶ vertex set V
- ▶ edge set $E \subseteq \left\{ \{u, v\} : u, v \in V \right\} \subset \mathcal{P}(V)$

First Betti number of a **topological space**:
maximum number of cuts that can be made
without dividing it into two separate pieces.

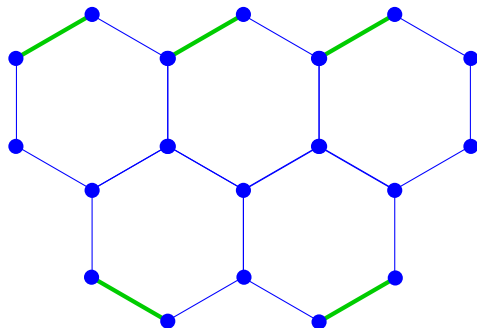
First Betti number of a **topological space**:
maximum number of cuts that can be made
without dividing it into two separate pieces.



Betti numbers / Informally

First Betti number of a **topological space**:

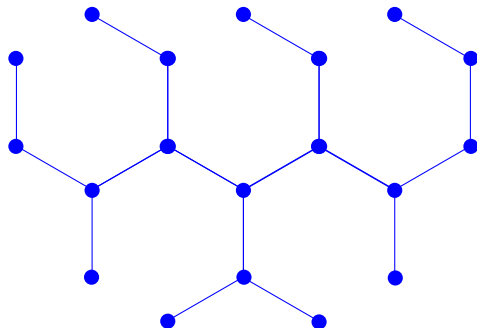
maximum number of cuts that can be made without dividing it into two separate pieces.



Betti numbers / Informally

First Betti number of a topological space:

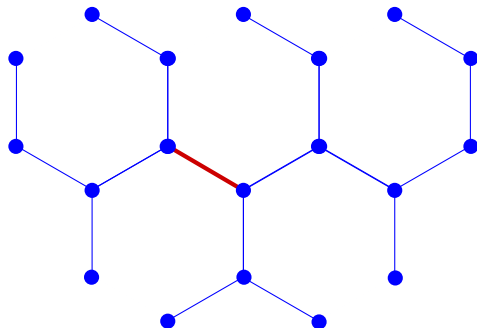
maximum number of cuts that can be made without dividing it into two separate pieces.



Betti numbers / Informally

First Betti number of a **topological space**:

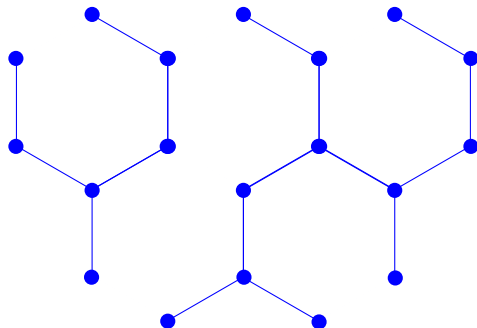
maximum number of cuts that can be made without dividing it into two separate pieces.



Betti numbers / Informally

First Betti number of a **topological space**:

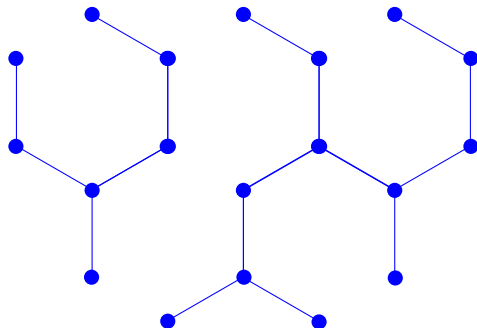
maximum number of cuts that can be made without dividing it into two separate pieces.



$$b_1 = 5$$

First Betti number of a **topological space**:

maximum number of cuts that can be made without dividing it into two separate pieces.



$$b_1 = 5$$

Geometric realization of G :
topological space where

Vertices $\longrightarrow$ points

Edges $\longrightarrow$ arcs that join vertices

Cycle Space

A **cycle** C in G is a closed path.

We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$

$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$

Cycle Space

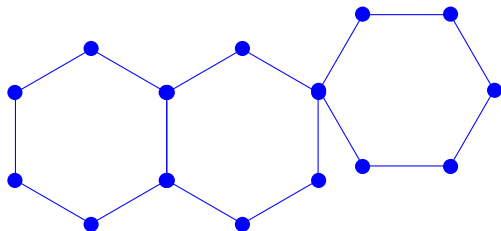
A **cycle** C in G is a closed path.

We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$
$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$



Cycle Space

A **cycle** C in G is a closed path.

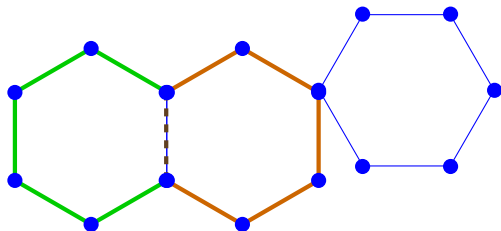
We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$

$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$



Cycle Space

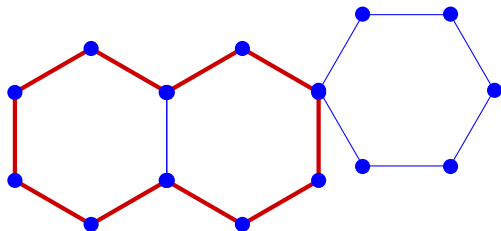
A **cycle** C in G is a closed path.

We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$
$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$



Cycle Space

A **cycle** C in G is a closed path.

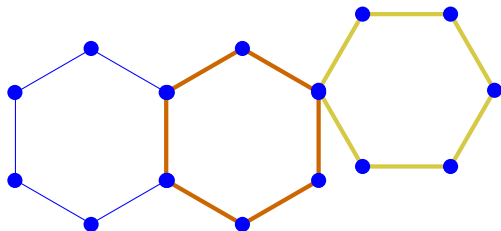
We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$

$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$



Cycle Space

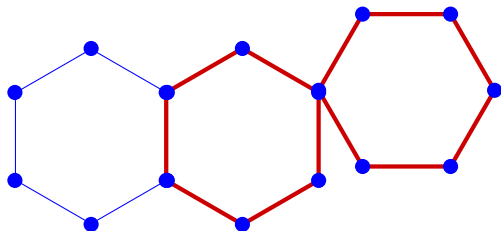
A **cycle** C in G is a closed path.

We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$
$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$



Cycle Space

A **cycle** C in G is a closed path.

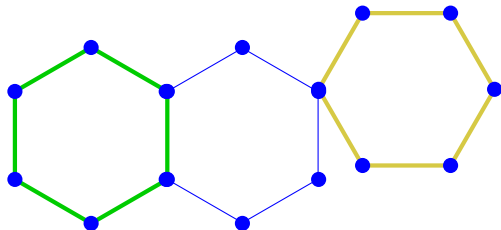
We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$

$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$



Cycle Space

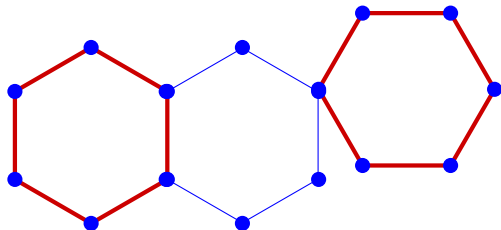
A **cycle** C in G is a closed path.

We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

$$S_1 \oplus S_2 = (S_1 \cup S_2) \setminus (S_1 \cap S_2)$$
$$1 \cdot S = S \quad \text{and} \quad 0 \cdot S = \emptyset$$



Cycle Space

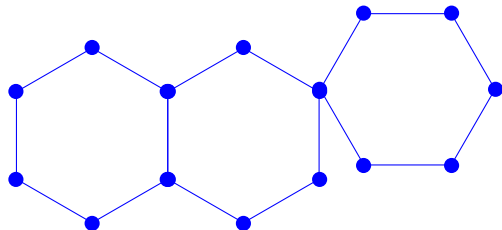
A **cycle** C in G is a closed path.

We identify C by its edge set, i.e., $C \in \mathcal{P}(E(G))$.

We can generalize cycles as arbitrary Eulerian subgraphs.

The set $\mathcal{C}(G)$ of all cycles forms a vector space over field $\mathbb{Z}_2$ with

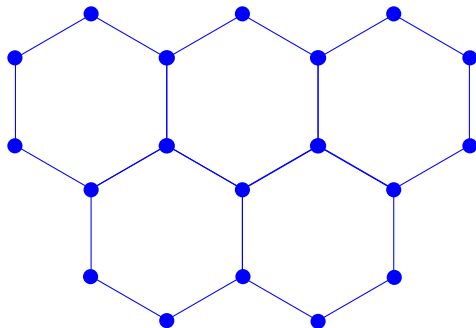
$$\begin{aligned} S_1 \oplus S_2 &= (S_1 \cup S_2) \setminus (S_1 \cap S_2) \\ 1 \cdot S &= S \quad \text{and} \quad 0 \cdot S = \emptyset \end{aligned}$$



$$\begin{aligned} \dim(\mathcal{C}) &= |E| - |V| + 1 \\ &= 17 - 15 + 1 = 3 \end{aligned}$$

cyclomatic number

Cycle Space



$$\begin{aligned}\dim(\mathcal{C}) &= |E| - |V| + 1 \\ &= 23 - 19 + 1 = 5\end{aligned}$$

CW complex

- (1) Start with a finite discrete set X^0 of 0-cells.
- (2) Inductively create n -skeleton X^n from X^{n-1} by attaching n -cells via continuous gluing maps $\partial_n D^n = S^{n-1} \rightarrow X^{n-1}$.

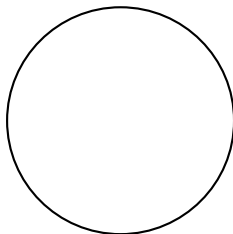
where

- ▶ **0-cell**: a single point
- ▶ **n -cell**: homeomorphic to open n -disc $D^n = \{x \in \mathbb{R}^n : \|x\| < 1\}$

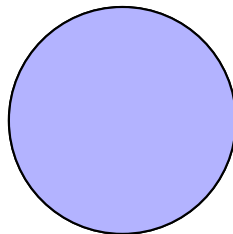
Skeleton X^n is a disjoint union of cells.

$$\emptyset = X^{-1} \subseteq X^0 \subseteq X^1 \subseteq X^2 \subseteq \dots \subseteq X$$

Circle:

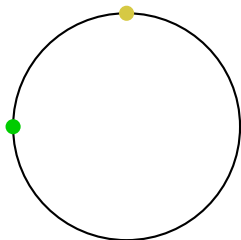


Disc:



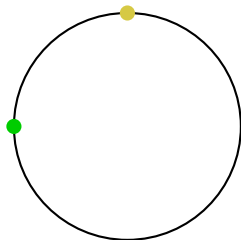
CW complex

Circle:



two **0-cells**

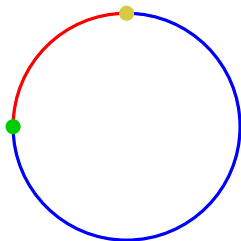
Disc:



two **0-cells**

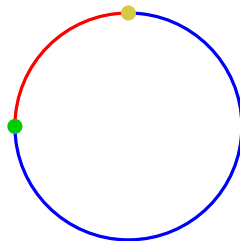
CW complex

Circle:



two **0-cells**
two **1-cells**

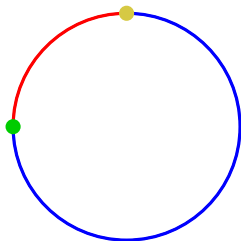
Disc:



two **0-cells**
two **1-cells**

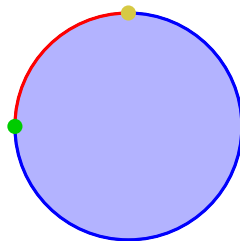
CW complex

Circle:



two **0-cells**
two **1-cells**

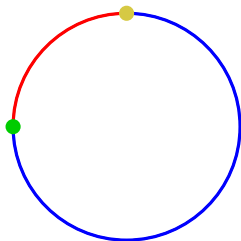
Disc:



two **0-cells**
two **1-cells**
one **2-cell**

CW complex

Circle:



1-dimensional CW complex

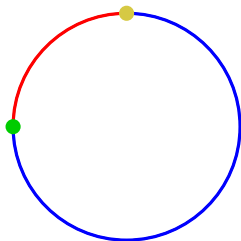
Disc:

A disc with a blue boundary. The interior of the disc is filled with light blue. A red arc is highlighted on the upper-left portion of the boundary. A green dot is located on the blue boundary at the start of the red arc, and a yellow dot is at the end of the red arc.

2-dimensional CW complex

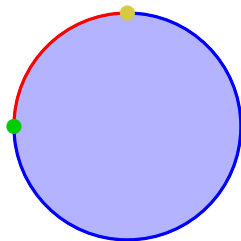
CW complex

Circle:



1-dimensional CW complex

Disc:



2-dimensional CW complex

The geometric realization of a graph is a 1-dimensional CW complex.

First Betti number of a **topological space**:

- ▶ **rank** of the first **homology group**.

n th Betti number of a **topological space**:

- ▶ **rank** of the n th **homology group**.

First Betti number of a **topological space**:

- ▶ **rank** of the first **homology group**.

n th Betti number of a **topological space**:

- ▶ **rank** of the n th **homology group**.

CW complexes are a convenient tool for computing homology groups
(and hence Betti numbers)

Chain complex

Let operator ∂_n assign the **boundary** for a n -dimensional region, e.g.,

$$\partial_3(D^3) = S^2 \quad \text{but} \quad \partial_3(D^2) = \emptyset$$

Also observe

$$\partial_3(D^3) = S^2 \quad \text{and} \quad \partial_2(\partial_3(D^3)) = \partial_2(S^2) = \emptyset$$

We get a chain

$$\dots \longrightarrow X^{n+1} \xrightarrow{\partial_{n+1}} X^n \xrightarrow{\partial_n} X^{n-1} \longrightarrow \dots$$

with

$$\partial_n \circ \partial_{n+1} = 0$$

Homology Group

The set of cells can be described by Abelian groups.
(Similar to cycle space.)

- ▶ cycle group

$$Z_n = \ker \partial_n$$

- ▶ boundary group

$$B_{n+1} = \operatorname{im} \partial_{n+1}$$

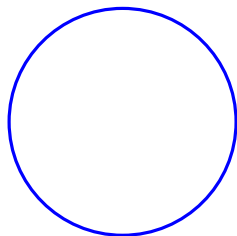
- ▶ n th homology group $H_n(X)$

$$H_n = Z_n / B_n = \ker \partial_n / \operatorname{im} \partial_{n+1}$$

Sets without boundary that are not the boundary of a set: **holes**

Homology Group

Circle:



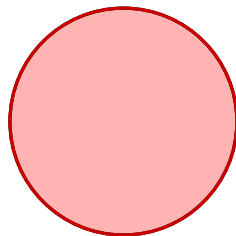
$$Z_1 = \mathbb{Z}$$

$$B_2 = \emptyset$$

$$H_1 = \mathbb{Z}/\emptyset = \mathbb{Z}$$

$$\dim(H_1) = 1$$

Disc:



$$Z_1 = \mathbb{Z}$$

$$B_2 = \mathbb{Z}$$

$$H_1 = \mathbb{Z}/\mathbb{Z} = 0$$

$$\dim(H_1) = 0$$

Boundary Operator on Graphs

Observe:

- ▶ 2-cells in geometric realizations can be identified by (connected) cycles.
- ▶ $\mathcal{P}(V)$ forms a vector space.
- ▶ Boundary of a union of edges is a subset of V .
- ▶ Cycle are exactly those with $\bigoplus_{c_i \in C} c_i = \emptyset$.

Boundary operator

$$\partial: \mathcal{P}(\mathcal{P}(V)) \rightarrow \mathcal{P}(V), \quad S \mapsto \partial S = \bigoplus_{s_i \in S} s_i$$

$$C \subseteq \mathcal{P}(V) \text{ is a cycle} \quad \Leftrightarrow \quad \partial C = \emptyset$$

Hypergraphs

A **simple** undirected **hypergraph** consists of

- ▶ **vertex** set V
- ▶ **hyperedges** $E \subseteq \mathcal{P}(V) \setminus \{\emptyset\}$

We have

- ▶ **boundary** operator

$$\partial: \mathcal{P}(\mathcal{P}(V)) \rightarrow \mathcal{P}(V), \quad S \mapsto \partial S = \bigoplus_{s_i \in S} s_i$$

- ▶ **cycle space**¹

$$\mathcal{C}(G) = \ker \partial \subseteq \mathcal{P}(E)$$

¹In literature a cycle in a hypergraph is defined as sequence of hyperedges

Matroid

Vector space $\mathcal{P}(E)$ forms a **matroid**:

- ▶ A is **linearly dependent** if it contains a cycle.
- ▶ **elementary cycle**: minimal linearly dependent set
Graph: a connected cycle where all degrees are 2.

Previous work by Döring (1975) and Oubiña (1974)

CW Complex on Hypergraph

Geometric realization of a graphs is a 1-dimensional CW complexes with

- ▶ vertex set V as its 0-skeleton and
- ▶ arcs as 1-cells,
- ▶ but no 2-cells. $\Rightarrow b_2(G) = 0$

CW Complex on Hypergraph

Geometric realization of a graphs is a 1-dimensional CW complexes with

- ▶ vertex set V as its 0-skeleton and
- ▶ arcs as 1-cells,
- ▶ but no 2-cells. $\Rightarrow b_2(G) = 0$

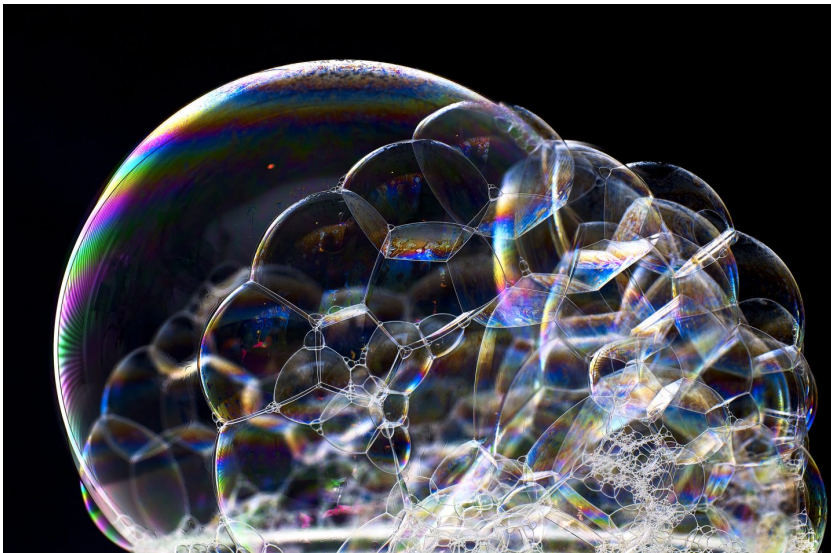
Idea:

- ▶ Attach 2-cells to geometric realization. $\Rightarrow X^2 \quad b_2(G)$
- ▶ Attach n -cells to geometric realization. $\Rightarrow X^n \quad b_n(G)$

Issue:

- ▶ Where do we attach n -cells?
- ▶ How can we describe this procedure?

Bubbles



Source: Getty Images

2-Cells on Graphs

- ▶ 0-cell: vertex in $V(G)$
- ▶ 1-cell: edge in $E(G)$
 - identified by its boundary
 - an element in $\mathcal{P}(V)$.
- ▶ 2-cell:
 - identified by its boundary
 - identified by (elementary) cycles:
an element in $\mathcal{C}(G) = \ker \partial \subseteq \mathcal{P}(E)$.

Select appropriate subset $E^2(G) \subseteq \mathcal{C}(G) \in \mathcal{P}(\mathcal{C}(G))$ (How?)

2-Cells on Graphs

- ▶ 0-cell: vertex in $V(G)$
- ▶ 1-cell: edge in $E(G)$
 - identified by its boundary
 - an element in $\mathcal{P}(V)$.
- ▶ 2-cell:
 - identified by its boundary
 - identified by (elementary) cycles:
an element in $\mathcal{C}(G) = \ker \partial \subseteq \mathcal{P}(E)$.

Select appropriate subset $E^2(G) \subseteq \mathcal{C}(G) \in \mathcal{P}(\mathcal{C}(G))$ (How?)

New hypergraph

$$G_1 = (E(G), E^2(G))$$

3-Cells on Graphs

- ▶ 3-cells: bubbles
 - identified by its boundary
 - identified by (elementary) cycles:
an element in $\ker \partial \subseteq \mathcal{P}(E^2)$.

Select appropriate subset $E^3(G) \subseteq \ker \partial \in \mathcal{P}(\ker \partial)$ (How?)

New hypergraph

$$G_2 = (E^2(G), E^3(G))$$

Edge Complex

$$\text{(C1)} \quad E^{-1} = \emptyset$$

$$\text{(C2)} \quad E^0 = V(G)$$

$$\text{(C3)} \quad E^1 = E(G) \subseteq \mathcal{P}(V(G)) = \ker \partial_0 = \mathcal{P}(E^0)$$

$$\text{(C4)} \quad E^{n+1} \subseteq \ker \partial_n \subseteq \mathcal{P}(E^n) \quad \text{for } n \geq 1,$$

$$\partial_0: \mathcal{P}(E^0) \rightarrow \mathcal{P}(E^{-1}) = \{\emptyset\}$$

$$\begin{aligned} \partial_n: \mathcal{P}(E^n) &\rightarrow \mathcal{P}(E^{n-1}), \\ \{e_1^n, \dots, e_l^n\} &\mapsto \partial(\{e_1^n, \dots, e_l^n\}) = \bigoplus_{i=1}^l e_i^n. \end{aligned}$$

Edge Complex

$$(C1) \quad E^{-1} = \emptyset$$

$$(C2) \quad E^0 = V(G)$$

$$(C3) \quad E^1 = E(G) \subseteq \mathcal{P}(V(G)) = \ker \partial_0 = \mathcal{P}(E^0)$$

$$(C4) \quad E^{n+1} \subseteq \ker \partial_n \subseteq \mathcal{P}(E^n) \quad \text{for } n \geq 1,$$

$$\partial_0: \mathcal{P}(E^0) \rightarrow \mathcal{P}(E^{-1}) = \{\emptyset\}$$

$$\begin{aligned} \partial_n: \mathcal{P}(E^n) &\rightarrow \mathcal{P}(E^{n-1}), \\ \{e_1^n, \dots, e_l^n\} &\mapsto \partial(\{e_1^n, \dots, e_l^n\}) = \bigoplus_{i=1}^l e_i^n. \end{aligned}$$

$$\begin{aligned} \dots \longrightarrow \mathcal{P}(E^{n+1}) &\xrightarrow{\partial_{n+1}} \mathcal{P}(E^n) \xrightarrow{\partial_n} \mathcal{P}(E^{n-1}) \longrightarrow \dots \\ \text{with } \partial_n \circ \partial_{n+1}(C) &= \emptyset \end{aligned}$$

Higher Order Betti Numbers

n th order cycle space:

$$\mathcal{C}_n(G) = \ker \partial_n \subseteq \mathcal{P}(E^n)$$

n th Betti number:

$$b_n(G) = \dim (\mathcal{C}_n(G))$$

Higher Order Betti Numbers

n th order cycle space:

$$\mathcal{C}_n(G) = \ker \partial_n \subseteq \mathcal{P}(E^n)$$

n th Betti number:

$$b_n(G) = \dim(\mathcal{C}_n(G))$$

Sequence of hypergraphs

$$G_n = (E^n, E^{n+1})$$

with

$$\mathcal{C}_n(G) = \mathcal{C}(G_n)$$

How? – Selector

We are in urgent need of a **selector** (**choice function**)

$$\mathcal{E}: G_{n-1} = (E^{n-1}, E^n) \mapsto E^{n+1} \subseteq \ker \partial_n \subseteq \mathcal{P}(E^n)$$

such that

(S1) $\mathcal{E}$ is universal, i.e., it does not depend on n and G .

(S2) All elements in the image of $\mathcal{E}$ are circuits in matroid $\mathcal{P}(E^n)$.

Recursion for G_n

$$G_n = (E^n, E^{n+1}) = (E^1(G_{n-1}), \mathcal{E}(G_{n-1}))$$

and

$$E^k(G_n) = E^{k+j}(G_{n-j}) = E^{k+n}(G)$$

n th cycle space $\mathcal{C}_n(G)$ heavily depends on selector $\mathcal{E}$.

So we have to use notations

$$\mathcal{C}_{n,\mathcal{E}}(G) \quad \text{and} \quad b_{n,\mathcal{E}}(G) = \dim(\mathcal{C}_{n,\mathcal{E}}(G))$$

and call $b_{n,\mathcal{E}}(G)$ the the n th Betti number w.r.t. selector $\mathcal{E}$.

Possible Selectors

- ▶ All **elementary cycles**:

$b_n(G)$ explodes for increasing n .

- ▶ Only those cycles that have strictly **larger support** than each of its edges:

$b_n(G) = 0$ for all $n \geq |V(G)|$ but is sensible against edge splitting

... and $b_k(G)$ can still be huge for $k < n$.

- ▶ Minimal length cycles:

Length is a metric concept and not a purely topological one.

Hence it is extremely sensible against edge splitting.

- ▶ Non-included cycles

- ▶ ???

The Big Question

- ▶ What do these Betti numbers **tell us** about a (hyper) graph?

Acknowledgments

I want to thank ...

Michael Hecht for raising this question.

Marc Hellmuth for lots of useful discussions.

Thank you for your attention!

References I

- B. Döring. Zyklen und Cozyklen in Hypergraphen. Diplomarbeit, Universität Bielefeld, 1975.
URL <http://www.bernd-doering.de/Zyklen%20und%20Cozyklen%20in%20Hypergraphen.pdf>.
- L. Oubiña. Des cycles et cocycles en hypergraphes. *Discrete Mathematics*, 8:71–80, 1974. doi: 10.1016/0012-365X(74)90111-3. <https://zbmath.org/0274.05119>.